# Remove imaginary types

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**target**

integration into IS ISO/IEC 9899:202y

**document history**

| document number | date | comment |
|---|---|---|
| n3240 (https://open-std.org/JTC1/SC22/WG14/www/docs/n3240.pdf) | **202404** | original proposal |
| n3263 (https://open-std.org/JTC1/SC22/WG14/www/docs/n3263.pdf) | **202404** | revive G.5, discuss `imaginary` and `_Generic` |
| n3274 (https://open-std.org/JTC1/SC22/WG14/www/docs/n3274.pdf) | **202406** | this document: take care of `I`, typo correction |

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# 1 Problem description

Optional imaginary types, indicated by the keyword `_Imaginary`, have no current implementation that would be known to WG14, and we are not aware of any plans for implementing them. The document n3206 (https://open-std.org/JTC1/SC22/WG14/www/docs/n3206.htm) described a number of problems with these types. It was discussed in the January 2024 meeting in Strasbourg and as a result there was consensus to remove this option from the C standard.

## 1.1 Possible existing implementations with imaginary types

Up to now we have no information about an up-to-date conforming implementation that would have imaginary types. Nonetheless, for most of it such a support would not make a hypothetical implementation behave badly for C2y. It would simply transition through imaginary types on rare occasions, but the user-visible results of complex type would be the same.

### 1.1.1 Reserved identifiers

The identifiers `_Imaginary` and `_Imaginary_I` are reserved, anyhow, so no conflict with these identifiers should occur.

The macros `imaginary` and `I` indeed step on the name space that is reserved for application use. But C23 has provisions to allow applications to undefine and redefine these macros, anyhow. So no conflict with user code should be expected. Nevertheless, we propose to add an explicit permissions to define `imaginary` and to use another implementation-defined value for `I`.

### 1.1.2 `_Generic` expressions using imaginary types

Hypothetical application code that uses `_Generic` with an association type that is imaginary will cease to work. Since there is probably no up-to-date compiler implementing imaginary types which would implement imaginary types and `_Generic` at the same time, the existence of such code in the field is even less likely.

# 2 Optional removal of `__STDC_IEC_559_COMPLEX__`

It seems that no recent implementation has correctly implemented this feature test macro which has been declared obsolete in C23. Therefor we propose to remove it together with the other changes. Such a removal is direct, so we do not explicitly propose wording below.

# 3 Questions

1. Does WG14 want to remove optional support for imaginary types as proposed in n3274 (https://open-std.org/JTC1/SC22/WG14/www/docs/n3274.pdf) from C2y?

2. Does WG14 want to remove the specification of `__STDC_IEC_559_COMPLEX__` from C2y?

3. Does WG14 want to apply the editorial changes as proposed in n3274 (https://open-std.org/JTC1/SC22/WG14/www/docs/n3274.pdf) for C2y?

# 4 Wording

Removals are in ~~stroke-out red~~, additions in <u>underlined green</u>.

## 4.1 Changes for `__STDC_IEC_559_COMPLEX__`

- Remove that macro from the predefined macros in 5.2.5.3.3, 6.10.10.4 and 6.11.7

## 4.2 Changes to Clause 6

- Remove the footnote from 6.2.5 (Types) p15

- 6.3.1.8 change the first indented item as follows:

  *If one operand has decimal floating type, the other operand shall not have standard floating~~, complex or imaginary~~ <u>or complex</u> type.*

- Remove `_Imaginary` from the list of keywords in 6.4.1 p1

- Remove the following sentence from 6.4.1 p2

  ~~*The keyword `_Imaginary` is reserved for specifying imaginary types.*~~

- Remove the footnote from 6.4.1 p2

- Remove the following phrase from 6.7.6 (Storage-class specifiers) p6

  ~~*If the object declared has imaginary type, the initializer shall have imaginary type.*~~

- Remove the following code snippet from 6.7.6 p21 (Example 5)

```
#ifdef __STDC_IEC_60559_COMPLEX__
constexpr double d5 = (double _Imaginary)0.0; // constraint violation
constexpr double d6 = (double _Imaginary)0.0; // constraint violation
constexpr double _Imaginary di1 = 0.0*I;      // ok
constexpr double _Imaginary di2 = 0.0;        // constraint violation
#endif
```

- Remove the mention of imaginary types from 6.5 (Expressions), 5 places.

## 4.3 Changes to Clause 7.3 (Complex arithmetic)

- Remove p6

- Change p7

  *The macro*

  ```
  I
  ```

  *expands to* ~~*either `_Imaginary_I` or*~~ *`_Complex_I`.* ~~*If `_Imaginary_I` is not defined, I shall expand to `_Complex_I`.*~~ <u>*or to another implementation-defined arithmetic constant expression that results in an imaginary unit when converted to a complex type.*</u>

- Change p8

  *Notwithstanding the provisions of 7.1.3,*

- *a program may undefine and perhaps then redefine the macros* `complex`*, ~~imaginary~~,, and* `I`*;*
- *the implementation may define a macro* `imaginary` *but a program may undefine and perhaps then redefine that identifier .*

- Remove the following sentence from the footnotes to 7.3.9.2 (The `cimag` functions) p2 and 7.3.9.6 (The `creal` functions) p2

  *~~If imaginary types are supported, z and creal(z)+cimag(z)*I are equivalent expressions.~~*

- In 7.3.9.3 (The `CMPLX` macros) remove p4, NOTE, explaining how these macros would work if there were such a thing as imaginary types.

## 4.4 Optional editorial change

With the above changes, the introductory text to 7.3 becomes a bit disorganized. We propose to reorganize it a bit with the following editorial change.

- Remove current paragraphs p6, p7 and p8.

- Include a new paragraph before current paragraph p4.

  *The macro* `complex` *expands to* `_Complex`*; the macro* `_Complex_I` *expands to an arithmetic constant expression of type* `float` `_Complex`*, with the value of the imaginary unit;[239] the macro* `I` *expands to* `_Complex_I` *or to another implementation-defined arithmetic constant expression that results in an imaginary unit when converted to a complex type. Notwithstanding the provisions of 7.1.13,*

  - *a program may undefine and perhaps then redefine the macros* `complex`*, and* `I`*;*
  - *the implementation may define a macro* `imaginary` *but a program may undefine and perhaps then redefine that identifier.*

And keep footnote [239]

  [239] *The imaginary unit is a number i such that* $i^2 = -1$*.*

## 4.5 Changes to Annex G

There are a lot of removals from Annex G, such that a diff is merely unreadable. An appendix to this document has the proposed replacement, supposing that `__STDC_IEC_559_COMPLEX__` is also removed from C2y.

# 5 Changes to the LaTeX source – notes to reviewers and editors

The branch "imaginary3" has the patches

- *6dd22a4* remove optional support of imaginary types from clauses 6 and 7
- *9d34847* take care of the naming issues concerning imaginary (new for n3274 (https://open-std.org/JTC1/SC22/WG14/www/docs/n3274.pdf) )
- *2f2df87* adapt G .5 (new for n3274 (https://open-std.org/JTC1/SC22/WG14/www/docs/n3274.pdf) )
- *859b9d7* remove support for imaginary types from annexes
- *5354dc1* remove `__STDC_IEC_559_COMPLEX__`
- *ea9dae2* editorial: reorganize the start of <complex.h> a bit

# 6 Appendix: replacement of Annex G

# Annex G
(normative)
## ISO/IEC 60559-compatible complex arithmetic

## G.1   Introduction

1   This annex supplements Annex F to specify complex arithmetic for compatibility with ISO/IEC 60559 real floating-point arithmetic. An implementation that defines `__STDC_IEC_60559_COMPLEX__` shall conform to the specifications in this annex.[444]

## G.2   Conventions

1   A complex value with at least one infinite part is regarded as an *infinity* (even if its other part is a quiet NaN). A complex value is a *finite number* if each of its parts is a finite number (neither infinite nor NaN). A complex value is a *zero* if each of its parts is a zero.

## G.3   Binary operators
## G.3.1   General

1   For most operand types, the value of the result of a binary operator with a complex operand is completely determined, with reference to real arithmetic, by the usual mathematical formula. For some operand types, the usual mathematical formula is problematic because of its treatment of infinities and because of undue overflow or underflow; in these cases the result satisfies certain properties (specified in G.3.2), but is not completely determined.

## G.3.2   Multiplicative operators

**Semantics**

1   If the operands are not both complex, then the result and floating-point exception behavior of the $*$ operator is defined by the usual mathematical formula:

| $*$ | $u$ | $u + iv$ |
|---|---|---|
| $x$ | $xu$ | $(xu) + i(xv)$ |
| $x + iy$ | $(xu) + i(yu)$ | |

2   If the second operand is not complex, then the result and floating-point exception behavior of the $/$ operator is defined by the usual mathematical formula:

| $/$ | $u$ |
|---|---|
| $x$ | $x/u$ |
| $x + iy$ | $(x/u) + i(y/u)$ |

3   The $*$ and $/$ operators satisfy the following infinity properties for all real and complex operands:[445]

   — if one operand is an infinity and the other operand is a nonzero finite number or an infinity, then the result of the $*$ operator is an infinity;

   — if the first operand is an infinity and the second operand is a finite number, then the result of the $/$ operator is an infinity;

   — if the first operand is a finite number and the second operand is an infinity, then the result of the $/$ operator is a zero;

   — if the first operand is a nonzero finite number or an infinity and the second operand is a zero, then the result of the $/$ operator is an infinity.

4   If both operands of the $*$ operator are complex or if the second operand of the $/$ operator is complex, the operator raises floating-point exceptions if appropriate for the calculation of the parts of the result, and may raise spurious floating-point exceptions.

---

[444]Implementations that do not define `__STDC_IEC_60559_COMPLEX__` are not required to conform to these specifications.
[445]These properties are already implied for those cases covered in the tables, but are required for all cases (at least where the state for `CX_LIMITED_RANGE` is "off").

5    EXAMPLE 1  Multiplication of **double _Complex** operands could be implemented as follows.

```c
#include <math.h>
#include <complex.h>

/* Multiply z * w ...*/
double complex _Cmultd(double complex z, double complex w)
{
      #pragma STDC FP_CONTRACT OFF
      double a, b, c, d, ac, bd, ad, bc, x, y;
      a = creal(z); b = cimag(z);
      c = creal(w); d = cimag(w);
      ac = a * c;   bd = b * d;
      ad = a * d;   bc = b * c;
      x = ac - bd;  y = ad + bc;
      if (isnan(x) && isnan(y)) {
            /* Recover infinities that computed as NaN+iNaN ... */
            int recalc = 0;
            if (isinf(a) || isinf(b)) { // z is infinite
                  /* "Box" the infinity and change NaNs in the other factor to 0 */
                  a = copysign(isinf(a) ? 1.0: 0.0, a);
                  b = copysign(isinf(b) ? 1.0: 0.0, b);
                  if (isnan(c)) c = copysign(0.0, c);
                  if (isnan(d)) d = copysign(0.0, d);
                  recalc = 1;
            }
            if (isinf(c) || isinf(d)) { // w is infinite
                  /* "Box" the infinity and change NaNs in the other factor to 0 */
                  c = copysign(isinf(c) ? 1.0: 0.0, c);
                  d = copysign(isinf(d) ? 1.0: 0.0, d);
                  if (isnan(a)) a = copysign(0.0, a);
                  if (isnan(b)) b = copysign(0.0, b);
                  recalc = 1;
            }
            if (!recalc && (isinf(ac) || isinf(bd) ||
                            isinf(ad) || isinf(bc))) {
                  /* Recover infinities from overflow by changing NaNs to 0 ... */
                  if (isnan(a)) a = copysign(0.0, a);
                  if (isnan(b)) b = copysign(0.0, b);
                  if (isnan(c)) c = copysign(0.0, c);
                  if (isnan(d)) d = copysign(0.0, d);
                  recalc = 1;
            }
            if (recalc) {
                  x = INFINITY * (a * c - b * d);
                  y = INFINITY * (a * d + b * c);
            }
      }
      return CMPLX(x, y);
}
```

6    This implementation achieves the required treatment of infinities at the cost of only one **isnan** test in ordinary (finite) cases. It is less than ideal in that undue overflow and underflow could occur.

7    EXAMPLE 2  Division of two **double _Complex** operands could be implemented as follows.

```c
#include <math.h>
#include <complex.h>

/* Divide z / w ... */
double complex _Cdivd(double complex z, double complex w)
{
```

```
          #pragma STDC FP_CONTRACT OFF
          double a, b, c, d, logbw, denom, x, y;
          int ilogbw = 0;
          a = creal(z); b = cimag(z);
          c = creal(w); d = cimag(w);
          logbw = logb(fmaximum_num(fabs(c), fabs(d)));
          if (isfinite(logbw)) {
                  ilogbw = (int)logbw;
                  c = scalbn(c, -ilogbw); d = scalbn(d, -ilogbw);
          }
          denom = c * c + d * d;
          x = scalbn((a * c + b * d) / denom, -ilogbw);
          y = scalbn((b * c - a * d) / denom, -ilogbw);

          /* Recover infinities and zeros that computed as NaN+iNaN;         */
          /* the only cases are nonzero/zero, infinite/finite, and finite/infinite, ...   */

          if (isnan(x) && isnan(y)) {
                  if ((denom == 0.0) &&
                          (!isnan(a) || !isnan(b))) {
                          x = copysign(INFINITY, c) * a;
                          y = copysign(INFINITY, c) * b;
                  }
                  else if ((isinf(a) || isinf(b)) &&
                          isfinite(c) && isfinite(d)) {
                          a = copysign(isinf(a) ? 1.0: 0.0, a);
                          b = copysign(isinf(b) ? 1.0: 0.0, b);
                          x = INFINITY * (a * c + b * d);
                          y = INFINITY * (b * c - a * d);
                  }
                  else if ((logbw == INFINITY) &&
                          isfinite(a) && isfinite(b)) {
                          c = copysign(isinf(c) ? 1.0: 0.0, c);
                          d = copysign(isinf(d) ? 1.0: 0.0, d);
                          x = 0.0 * (a * c + b * d);
                          y = 0.0 * (b * c - a * d);
                  }
          }
          return CMPLX(x, y);
  }
```

8   Scaling the denominator alleviates the main overflow and underflow problem, which is more serious than for multiplication. In the spirit of the preceding multiplication example, this code does not defend against overflow and underflow in the calculation of the numerator. Scaling with the **scalbn** function, instead of with division, provides better roundoff characteristics.

## G.3.3   Additive operators

**Semantics**

1   In all cases the result and floating-point exception behavior of a + or - operator is defined by the usual mathematical formula:

| + or - | $u$ | $u + iv$ |
|---|---|---|
| $x$ | $x \pm u$ | $(x \pm u) \pm iv$ |
| $x + iy$ | $(x \pm u) + iy$ | $(x \pm u) + i(y \pm v)$ |

## G.4   Complex arithmetic **<complex.h>**

## G.4.1   General

1   This subclause contains specifications for the <complex.h> functions that are particularly suited to ISO/IEC 60559 implementations. For families of functions, the specifications apply to all of the

functions even though only the principal function is shown. Unless otherwise specified, where the symbol "$\pm$" occurs in both an argument and the result, the result has the same sign as the argument.

2  The functions are continuous onto both sides of their branch cuts, taking into account the sign of zero. For example, **csqrt**$(-2\pm i0) = \pm i\sqrt{2}$.

3  Since complex values are composed of real values, each function may be regarded as computing real values from real values. Except as noted, the functions treat real infinities, NaNs, signed zeros, subnormals, and the floating-point exception flags in a manner consistent with the specifications for real functions in F.10.[446]

4  In subsequent subclauses in G.4 "NaN" refers to a quiet NaN. The behavior of signaling NaNs in this annex is implementation-defined.

5  The functions **cimag**, **conj**, **cproj**, and **creal** are fully specified for all implementations, including ISO/IEC 60559 ones, in 7.3.9. These functions raise no floating-point exceptions.

6  Each of the functions **cabs** and **carg** is specified by a formula in terms of a real function (whose special cases are covered in Annex F):

```
cabs(x + iy)  =  hypot(x, y)
carg(x + iy)  =  atan2(y, x)
```

7  Each of the functions **casin**, **catan**, **ccos**, **csin**, and **ctan** is specified implicitly by a formula in terms of other complex functions (whose special cases are specified below):

```
casin(z)  =  −i casinh(iz)
catan(z)  =  −i catanh(iz)
ccos(z)   =  ccosh(iz)
csin(z)   =  −i csinh(iz)
ctan(z)   =  −i ctanh(iz)
```

8  For the other functions, the following subclauses specify behavior for special cases, including treatment of the "invalid" and "divide-by-zero" floating-point exceptions. For families of functions, the specifications apply to all of the functions even though only the principal function is shown. For a function $f$ satisfying $f(\mathrm{conj}(z)) = \mathrm{conj}(f(z))$, the specifications for the upper half-plane imply the specifications for the lower half-plane; if the function $f$ is also either even, $f(-z) = f(z)$, or odd, $f(-z) = -f(z)$, then the specifications for the first quadrant imply the specifications for the other three quadrants.

9  In the following subclauses, $\mathrm{cis}(y)$ is defined as $\cos(y) + i\sin(y)$.

## G.4.2  Trigonometric functions

### G.4.2.1  The **cacos** functions

1  — **cacos**$(\mathrm{conj}(z)) = \mathrm{conj}(\mathbf{cacos}(z))$.

— **cacos**$(\pm 0 + i0)$ returns $\frac{\pi}{2} - i0$.

— **cacos**$(\pm 0 + i\mathrm{NaN})$ returns $\frac{\pi}{2} + i\mathrm{NaN}$.

— **cacos**$(x + i\infty)$ returns $\frac{\pi}{2} - i\infty$, for finite $x$.

— **cacos**$(x + i\mathrm{NaN})$ returns $\mathrm{NaN} + i\mathrm{NaN}$ and optionally raises the "invalid" floating-point exception, for nonzero finite $x$.

— **cacos**$(-\infty + iy)$ returns $\pi - i\infty$, for positive-signed finite $y$.

— **cacos**$(+\infty + iy)$ returns $+0 - i\infty$, for positive-signed finite $y$.

— **cacos**$(-\infty + i\infty)$ returns $3\frac{\pi}{4} - i\infty$.

---

[446]As noted in G.2, a complex value with at least one infinite part is regarded as an infinity even if its other part is a quiet NaN.

— **cacos**$(+\infty + i\infty)$ returns $\frac{\pi}{4} - i\infty$.

— **cacos**$(\pm\infty + i\text{NaN})$ returns $\text{NaN}\pm i\infty$ (where the sign of the imaginary part of the result is unspecified).

— **cacos**$(\text{NaN} + iy)$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite $y$.

— **cacos**$(\text{NaN} + i\infty)$ returns $\text{NaN} - i\infty$.

— **cacos**$(\text{NaN} + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$.

## G.4.3 Hyperbolic functions

### G.4.3.1 The **cacosh** functions

1 — **cacosh**(**conj**$(z)$) = **conj**(**cacosh**$(z)$).

— **cacosh**$(\pm 0 + i0)$ returns $+0 + \frac{i\pi}{2}$.

— **cacosh**$(x + i\infty)$ returns $+\infty + \frac{i\pi}{2}$, for finite $x$.

— **cacosh**$(0 + i\text{NaN})$ returns $\text{NaN}\pm\frac{i\pi}{2}$ (where the sign of the imaginary part of the result is unspecified).

— **cacosh**$(x + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite nonzero $x$.

— **cacosh**$(-\infty + iy)$ returns $+\infty + i\pi$, for positive-signed finite $y$.

— **cacosh**$(+\infty + iy)$ returns $+\infty + i0$, for positive-signed finite $y$.

— **cacosh**$(-\infty + i\infty)$ returns $+\infty + i\frac{3\pi}{4}$.

— **cacosh**$(+\infty + i\infty)$ returns $+\infty + \frac{i\pi}{4}$.

— **cacosh**$(\pm\infty + i\text{NaN})$ returns $+\infty + i\text{NaN}$.

— **cacosh**$(\text{NaN} + iy)$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite $y$.

— **cacosh**$(\text{NaN} + i\infty)$ returns $+\infty + i\text{NaN}$.

— **cacosh**$(\text{NaN} + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$.

### G.4.3.2 The **casinh** functions

1 — **casinh**(**conj**$(z)$) = **conj**(**casinh**$(z)$). and **casinh** is odd.

— **casinh**$(+0 + i0)$ returns $0 + i0$.

— **casinh**$(x + i\infty)$ returns $+\infty + \frac{i\pi}{2}$ for positive-signed finite $x$.

— **casinh**$(x + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite $x$.

— **casinh**$(+\infty + iy)$ returns $+\infty + i0$ for positive-signed finite $y$.

— **casinh**$(+\infty + i\infty)$ returns $+\infty + \frac{i\pi}{4}$.

— **casinh**$(+\infty + i\text{NaN})$ returns $+\infty + i\text{NaN}$.

— **casinh**$(\text{NaN} + i0)$ returns $\text{NaN} + i0$.

— **casinh**$(\text{NaN} + iy)$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite nonzero $y$.

— **casinh**$(\text{NaN} + i\infty)$ returns $\pm\infty + i\text{NaN}$ (where the sign of the real part of the result is unspecified).

— **casinh**$(\text{NaN} + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$.

### G.4.3.3 The `catanh` functions

1
— `catanh(conj($z$)) = conj(catanh($z$))`. and `catanh` is odd.

— `catanh($+0 + i0$)` returns $+0 + i0$.

— `catanh($+0 + i$NaN)` returns $+0 + i$NaN.

— `catanh($+1 + i0$)` returns $+\infty + i0$ and raises the "divide-by-zero" floating-point exception.

— `catanh($x + i\infty$)` returns $+0 + \frac{i\pi}{2}$, for finite positive-signed $x$.

— `catanh($x + i$NaN)` returns NaN $+ i$NaN and optionally raises the "invalid" floating-point exception, for nonzero finite $x$.

— `catanh($+\infty + iy$)` returns $+0 + \frac{i\pi}{2}$, for finite positive-signed $y$.

— `catanh($+\infty + i\infty$)` returns $+0 + \frac{i\pi}{2}$.

— `catanh($+\infty + i$NaN)` returns $+0 + i$NaN.

— `catanh(NaN $+ iy$)` returns NaN $+ i$NaN and optionally raises the "invalid" floating-point exception, for finite $y$.

— `catanh(NaN $+ i\infty$)` returns $\pm 0 + \frac{i\pi}{2}$ (where the sign of the real part of the result is unspecified).

— `catanh(NaN $+ i$NaN)` returns NaN $+ i$NaN.

### G.4.3.4 The `ccosh` functions

1
— `ccosh(conj($z$)) = conj(ccosh($z$))` and `ccosh` is even.

— `ccosh($+0 + i0$)` returns $1 + i0$.

— `ccosh($+0 + i\infty$)` returns NaN$\pm i0$ (where the sign of the imaginary part of the result is unspecified) and raises the "invalid" floating-point exception.

— `ccosh($+0 + i$NaN)` returns NaN$\pm i0$ (where the sign of the imaginary part of the result is unspecified).

— `ccosh($x + i\infty$)` returns NaN $+ i$NaN and raises the "invalid" floating-point exception, for finite nonzero $x$.

— `ccosh($x + i$NaN)` returns NaN $+ i$NaN and optionally raises the "invalid" floating-point exception, for finite nonzero $x$.

— `ccosh($+\infty + i0$)` returns $+\infty + i0$.

— `ccosh($+\infty + iy$)` returns $+\infty$ cis($y$), for finite nonzero $y$.

— `ccosh($+\infty + i\infty$)` returns $\pm\infty + i$NaN (where the sign of the real part of the result is unspecified) and raises the "invalid" floating-point exception.

— `ccosh($+\infty + i$NaN)` returns $+\infty + i$NaN.

— `ccosh(NaN $+ i0$)` returns NaN$\pm i0$ (where the sign of the imaginary part of the result is unspecified).

— `ccosh(NaN $+ iy$)` returns NaN $+ i$NaN and optionally raises the "invalid" floating-point exception, for all nonzero numbers $y$.

— `ccosh(NaN $+ i$NaN)` returns NaN $+ i$NaN.

### G.4.3.5 The `csinh` functions

1    — `csinh`(`conj`($z$)) = `conj`(`csinh`($z$)). and `csinh` is odd.

   — `csinh`($+0 + i0$) returns $+0 + i0$.

   — `csinh`($+0 + i\infty$) returns $\pm 0 + i$NaN (where the sign of the real part of the result is unspecified) and raises the "invalid" floating-point exception.

   — `csinh`($+0 + i$NaN) returns $\pm 0 + i$NaN (where the sign of the real part of the result is unspecified).

   — `csinh`($x + i\infty$) returns NaN $+ i$NaN and raises the "invalid" floating-point exception, for positive finite $x$.

   — `csinh`($x + i$NaN) returns NaN $+ i$NaN and optionally raises the "invalid" floating-point exception, for finite nonzero $x$.

   — `csinh`($+\infty + i0$) returns $+\infty + i0$.

   — `csinh`($+\infty + iy$) returns $+\infty \; \mathrm{cis}(y)$, for positive finite $y$.

   — `csinh`($+\infty + i\infty$) returns $\pm\infty + i$NaN (where the sign of the real part of the result is unspecified) and raises the "invalid" floating-point exception.

   — `csinh`($+\infty + i$NaN) returns $\pm\infty + i$NaN (where the sign of the real part of the result is unspecified).

   — `csinh`(NaN $+ i0$) returns NaN $+ i0$.

   — `csinh`(NaN $+ iy$) returns NaN $+ i$NaN and optionally raises the "invalid" floating-point exception, for all nonzero numbers $y$.

   — `csinh`(NaN $+ i$NaN) returns NaN $+ i$NaN.

### G.4.3.6 The `ctanh` functions

1    — `ctanh`(`conj`($z$)) = `conj`(`ctanh`($z$)) and `ctanh` is odd.

   — `ctanh`($+0 + i0$) returns $+0 + i0$.

   — `ctanh`($0 + i\infty$) returns $0 + i$NaN and raises the "invalid" floating-point exception.

   — `ctanh`($x + i\infty$) returns NaN $+ i$NaN and raises the "invalid" floating-point exception, for finite nonzero $x$.

   — `ctanh`($0 + i$NaN) returns $0 + i$NaN.

   — `ctanh`($x + i$NaN) returns NaN $+ i$NaN and optionally raises the "invalid" floating-point exception, for finite nonzero $x$.

   — `ctanh`($+\infty + iy$) returns $1 + i0\sin(2y)$, for positive-signed finite $y$.

   — `ctanh`($+\infty + i\infty$) returns $1 \pm i0$ (where the sign of the imaginary part of the result is unspecified).

   — `ctanh`($+\infty + i$NaN) returns $1 \pm i0$ (where the sign of the imaginary part of the result is unspecified).

   — `ctanh`(NaN $+ i0$) returns NaN $+ i0$.

   — `ctanh`(NaN $+ iy$) returns NaN $+ i$NaN and optionally raises the "invalid" floating-point exception, for all nonzero numbers $y$.

   — `ctanh`(NaN $+ i$NaN) returns NaN $+ i$NaN.

### G.4.4  Exponential and logarithmic functions

#### G.4.4.1  The `cexp` functions

1
— $\text{cexp}(\text{conj}(z)) = \text{conj}(\text{cexp}(z))$.

— $\text{cexp}(\pm 0 + i0)$ returns $1 + i0$.

— $\text{cexp}(x + i\infty)$ returns $\text{NaN} + i\text{NaN}$ and raises the "invalid" floating-point exception, for finite $x$.

— $\text{cexp}(x + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite $x$.

— $\text{cexp}(+\infty + i0)$ returns $+\infty + i0$.

— $\text{cexp}(-\infty + iy)$ returns $+0\ \text{cis}(y)$, for finite $y$.

— $\text{cexp}(+\infty + iy)$ returns $+\infty\ \text{cis}(y)$, for finite nonzero $y$.

— $\text{cexp}(-\infty + i\infty)$ returns $\pm 0 \pm i0$ (where the signs of the real and imaginary parts of the result are unspecified).

— $\text{cexp}(+\infty + i\infty)$ returns $\pm\infty + i\text{NaN}$ and raises the "invalid" floating-point exception (where the sign of the real part of the result is unspecified).

— $\text{cexp}(-\infty + i\text{NaN})$ returns $\pm 0 \pm i0$ (where the signs of the real and imaginary parts of the result are unspecified).

— $\text{cexp}(+\infty + i\text{NaN})$ returns $\pm\infty + i\text{NaN}$ (where the sign of the real part of the result is unspecified).

— $\text{cexp}(\text{NaN} + i0)$ returns $\text{NaN} + i0$.

— $\text{cexp}(\text{NaN} + iy)$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for all nonzero numbers $y$.

— $\text{cexp}(\text{NaN} + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$.

#### G.4.4.2  The `clog` functions

1
— $\text{clog}(\text{conj}(z)) = \text{conj}(\text{clog}(z))$.

— $\text{clog}(-0 + i0)$ returns $-\infty + i\pi$ and raises the "divide-by-zero" floating-point exception.

— $\text{clog}(+0 + i0)$ returns $-\infty + i0$ and raises the "divide-by-zero" floating-point exception.

— $\text{clog}(x + i\infty)$ returns $+\infty + \frac{i\pi}{2}$, for finite $x$.

— $\text{clog}(x + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite $x$.

— $\text{clog}(-\infty + iy)$ returns $+\infty + i\pi$, for finite positive-signed $y$.

— $\text{clog}(+\infty + iy)$ returns $+\infty + i0$, for finite positive-signed $y$.

— $\text{clog}(-\infty + i\infty)$ returns $+\infty + i\frac{3\pi}{4}$.

— $\text{clog}(+\infty + i\infty)$ returns $+\infty + \frac{i\pi}{4}$.

— $\text{clog}(\pm\infty + i\text{NaN})$ returns $+\infty + i\text{NaN}$.

— $\text{clog}(\text{NaN} + iy)$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite $y$.

— $\text{clog}(\text{NaN} + i\infty)$ returns $+\infty + i\text{NaN}$.

— $\text{clog}(\text{NaN} + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$.

### G.4.5   Power and absolute-value functions

#### G.4.5.1   The **cpow** functions

1   The **cpow** functions raise floating-point exceptions if appropriate for the calculation of the parts of the result, and may also raise spurious floating-point exceptions.[447]

#### G.4.5.2   The **csqrt** functions

1   — **csqrt**(**conj**$(z)$) = **conj**(**csqrt**$(z)$).

— **csqrt**$(\pm 0 + i0)$ returns $+0 + i0$.

— **csqrt**$(x + i\infty)$ returns $+\infty + i\infty$, for all $x$ (including NaN).

— **csqrt**$(x + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite $x$.

— **csqrt**$(-\infty + iy)$ returns $+0 + i\infty$, for finite positive-signed $y$.

— **csqrt**$(+\infty + iy)$ returns $+\infty + i0$, for finite positive-signed $y$.

— **csqrt**$(-\infty + i\text{NaN})$ returns $\text{NaN}\pm i\infty$ (where the sign of the imaginary part of the result is unspecified).

— **csqrt**$(+\infty + i\text{NaN})$ returns $+\infty + i\text{NaN}$.

— **csqrt**$(\text{NaN} + iy)$ returns $\text{NaN} + i\text{NaN}$ and optionally raises the "invalid" floating-point exception, for finite $y$.

— **csqrt**$(\text{NaN} + i\text{NaN})$ returns $\text{NaN} + i\text{NaN}$.

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[447]This allows **cpow**$(z, c)$ to be implemented as **cexp**$(c\,\text{\textbf{clog}}(z))$ without precluding implementations that treat special cases more carefully.