

1 ISO/IEC JTC 1/SC 22/WG 23 N 0305

2 *Proposed separation of XYY into two descriptions*

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Date 15 February 2011
Contributed by Jim Moore
Original file name
Notes Responds to Action Item #16-12

The text of Action Item #16-12 reads as follows:

Look at XYY in the main document and both annexes to try to tease apart two vulnerabilities: one concerning arithmetic over/underflow and one concerning performing bit/shift operations on numeric values. In both, note that unsigned and signed arithmetic present two different challenges.

After reading the annexes, I realized that redrafting them will be easy once we settle on the text for the body of the report. So, I'm not including text for annexes at this time.

The proposed text for the body of the report follows:

6.x Arithmetic Wrap-around Error [FIF]

6.x.1 Description of application vulnerability

Wrap-around errors can occur whenever a value is incremented past the maximum or decremented past the minimum value representable in its type and, depending upon:

- whether the type is signed or unsigned
- the specification of the language semantics and/or
- implementation choices,

"wraps around" to an unexpected value. This vulnerability is related to Logical Wrap-around Error [PIK]. This description is derived from Wrap-Around Error [XYY], which appeared in Edition 1 of this international technical report.

6.x.2 Cross reference [Note to editor: Please verify the applicability of these cross-references.]

CWE:

- 128. Wrap-around Error
- 190: Integer Overflow or Wraparound

JSF AV Rules: 164 and 15

MISRA C 2004: 10.1 to 10.6, 12.8 and 12.11

MISRA C++ 2008: 2-13-3, 5-0-3 to 5-0-10, and 5-19-1

CERT C guidelines: INT30-C, INT32-C, and INT34-C

6.x.3 Mechanism of failure

Due to how arithmetic is performed by computers, if a variable's value is increased past the maximum value representable in its type, the system may fail to provide an overflow indication

42 to the program. One of the most common processor behaviour is to “wrap” to a very large
43 negative value, or set a condition flag for overflow or underflow, or saturate at the largest
44 representable value.

45
46 Wrap-around often generates an unexpected negative value; this unexpected value may cause a
47 loop to continue for a long time (because the termination condition requires a value greater than
48 some positive value) or an array bounds violation. A wrap-around can sometimes trigger buffer
49 overflows that can be used to execute arbitrary code.

50
51 It should be noted that the precise consequences of wrap-around differ depending on:

- 52 • Whether the type is signed or unsigned
- 53 • Whether the type is a modulus type
- 54 • Whether the type’s range is violated by exceeding the maximum representable value or
55 falling short of the minimum representable value
- 56 • The semantics of the language specification
- 57 • Implementation decisions

58 However, in all cases, the resulting problem is that the value yielded by the computation may be
59 unexpected.

60
61 6.x.4 Applicable language characteristics

62
63 This vulnerability description is intended to be applicable to languages with the following
64 characteristics:

- 65 • Languages that do not trigger an exception condition when a wrap-around error occurs.

66
67 6.x.4 Avoiding the vulnerability or mitigating its effects

68
69 Software developers can avoid the vulnerability or mitigate its ill effects in the following ways:

- 70 • Determine applicable upper and lower bounds for the range of all variables and use
71 language mechanisms or static analysis to determine that values are confined to the
72 proper range.
- 73 • Analyze the software using static analysis looking for unexpected consequences of
74 arithmetic operations.

75
76 6.x.6 Implications for standardization

77
78 In future standardization activities, the following items should be considered:

- 79 • Language standards developers should consider providing facilities to specify either an
80 error, a saturated value, or a modulo result when numeric overflow occurs. Ideally, the
81 selection among these alternatives could be made by the programmer.

82
83 6.y Logical Wrap-around Error [PIK]

84
85 6.y.1 Description of application vulnerability

86

87 Using shift operations as a surrogate for multiply or divide may produce an unexpected value
88 when significant bits are lost. This vulnerability is related to Arithmetic Wrap-around Error
89 [FIF]. This description is derived from Wrap-Around Error [XYY], which appeared in Edition 1
90 of this international technical report.

91
92 6.x.2 Cross reference [Note to editor: Please verify the applicability of these items.]
93

94 CWE:

95 128. Wrap-around Error
96 190: Integer Overflow or Wraparound
97 JSF AV Rules: 164 and 15
98 MISRA C 2004: 10.1 to 10.6, 12.8 and 12.11
99 MISRA C++ 2008: 2-13-3, 5-0-3 to 5-0-10, and 5-19-1
100 CERT C guidelines: INT30-C, INT32-C, and INT34-C

101
102 6.y.3 Mechanism of failure
103

104 Coders sometimes use shift operations with the intention of producing results equivalent to
105 multiplying by a power of two or dividing by a power of two. However, errors can result from
106 this practice. For example, if the programmer mistakenly uses logical shifts on signed arithmetic
107 values, the results may test correctly for small values but produce unexpected results when used
108 with large values. The problem, of course, is that the sign bit can be shifted out of the value
109 converting a negative value into a positive one or vice versa.

110
111 Even when the correct type of shift is coded, there can still be problems with unexpected and
112 undetected numerical underflow or overflow if significant bits are shifted out of the value.

113
114 Stated most generally, replacing multiply and divide operations with shifting operations requires
115 detailed knowledge of the representation of the values across the varieties of processors on which
116 the code may be used. In addition, it requires detailed analysis of the range of values for which
117 the shift operations will produce valid results and checking (or static analysis) to ensure that the
118 values never go outside of the range.

119
120 Wrap-around often generates an unexpected negative value; this unexpected value may cause a
121 loop to continue for a long time (because the termination condition requires a value greater than
122 some positive value) or an array bounds violation. A wrap-around can sometimes trigger buffer
123 overflows that can be used to execute arbitrary code.

124
125 6.y.4 Applicable language characteristics
126

127 This vulnerability description is intended to be applicable to languages with the following
128 characteristics:

- 129 • Languages that do not trigger an exception condition when a wrap-around error occurs.
- 130 • Languages that do not fully specify the distinction between arithmetic and logical shifts.

131
132 6.y.4 Avoiding the vulnerability or mitigating its effects

133

134 Software developers can avoid the vulnerability or mitigate its ill effects in the following ways:

- 135 • Determine applicable upper and lower bounds for the range of all variables and use
136 language mechanisms or static analysis to determine that values are confined to the
137 proper range.
- 138 • Analyze the software using static analysis looking for unexpected consequences of shift
139 operations.
- 140 • Avoid using shift operations as a surrogate for multiplication and division. Most
141 compilers will use the correct operation in the appropriate fashion when it is applicable.

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143 6.y.6 Implications for standardization

144

145 In future standardization activities, the following items should be considered:

- 146 • Language standards developers should consider providing facilities to specify either an
147 error, a saturated value, or a modulo result when logical overflow occurs. Ideally, the
148 selection among these alternatives could be made by the programmer.

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