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24772-1 DIS Collated comments with proposed edits

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| CA 001 |  | 06.02.5 |  | ed | "Avoid explicit type conversion of data values except when there is no alternative. Document suchoccurrences so that the justification is made available to maintainers."This could be read as encouraging implicit type conversions over explicit ones, which Ibelieve is the opposite of the intent. Further, while I can possibly imagine it might makesense to document implicit conversions, I think the explicit conversion mostlyreduces or eliminates the need to document the conversions. There is no mention in the previous bulletabout documenting implicit conversions, where the need is greater. | Suggest removing "explicit" from this bullet, i.e."Avoid [explicit] type conversion of data values except when there is no alternative. Document suchoccurrences so that the justification is made available to maintainers." | To be discussed |
| CA 002 |  | 06.03.3 |  | ed | "Problems can arise when programmers mix their techniques to reference the bits or output the bits. Problems can arisewhen programmers mix arithmetic and logical operations to reference the bits or output the bits. " There appears to be two versions of this sentence. Are they both needed?  | Delete one of the two sentences if both are not needed. Otherwise, the second sentence should probably begin with "Problems can {also} arise ...." | Accept.Combined intoProblems can arise when programmers mix their techniques to reference the bits or output the bits; or mix arithmetic and logical operations to reference the bits or output the bits. |
| PL 003 |  | 06.04 | 06.4.5 | te | Last subpoint: should be completed by "or use a suitable compensated summation algorithm" to read | When adding (or subtracting) sequences of numbers, sort and add (or subtract) them from smallest to largest in absolute value or use a suitable compensated summation algorithm to avoid loss of precision. | Accept. Added suggested text. |
| CA 004 |  | 06.08.3 |  | ed | “When an array has been allocated storage on the stack{,} an out-of-bounds write access may modifyinternal runtime housekeeping information (for example, a function's return address) which might changea program’s control flow” | Add missing comma | Fixed. |
| CA 005 |  | 06.14.06 |  | te | "Language specifiers should design generics in such a way that any attempt to instantiate a generic withconstructs that do not provide the required capabilities results in a compile-time error." | This seems completely unrelated to the vulnerability, nothing to do with free, but does sound like it could be applied to a different vulnerability relating to generics and templates. Is this out of place? Perhaps a cut and paste error? | Agreed. removed. |
| CA 006 |  | 06.17.04 |  | te | "Languages that treat letter case as significant. Some languages do not differentiate between names withdiffering case, while others do." This suggests that the problem is with languages that treat case as significant, but the description in 6.17.1, in the 4th paragraph states; “There is also an issue where identifiers appear distinct to a human but identical to the computer”implies the problem is with languages where case is insignificant. In reality, I think confusion can comefrom either type of language. But mostly only in cases of languages where variables do not need to be declared. | Suggest replacing the 4th bullet with a more critical issue;“Languages that allow variables to be implicitly declared.”An example that comes to mind here is BASIC, where, at least in earlier versions of the language, one could misspell the name of a variable and implicitly get a second variable, which might not be noticed by the programmer. | Accept in principle. The explanation in 6.17.1 was incomplete. Added “or the reverse where the human considers FOO and foo to be the same but the language treats them as distinct” to help explain the guidance in 6,17.4.The suggested additional guidance will be considered in a future revision. |
| CA 007 |  | 06.22.03 |  | te | This vulnerability appears to cover elaboration errors? That is another source of such an error that is not described here, where a global variable can be used before initialization,even though in the code, the variable appears to be initialized. I do see it mentioned in 6.22.5. | This can be a very dangerous problem as the code looks correct. I think it would be helpful to describe this possibility better in 6.22.1 and 6.22.3 | Reject for this document. The access to a variable during elaboration (an Ada concept) but with a later initialization is covered by this general description.  |
| CA 008 |  | 06.22.03 |  | ed | "There is a special case [of]{for} pointers or access types." | Replace “of” with “for”, or perhaps “associated with” | Accept.  |
| CA 009 | 4th bullet | 06.22.05 |  | te | There is a recommendation to initializing variables during elaboration, but that is precisely what causes this problem in most languages such as C++. The recommendation to initialize during elaboration only is good advice for languages that can perform safe/orderly elaboration based on dependencies, so there should be clarification. A general recommendation for all languages that allow global declarations would be to avoid creating global objects. For example, dont create a singleton object if all that is needed is function calls that generate a result based on the input parameters. For languages that do not support safe elaboration, the recommendation should be to NOT initialize the object during elaboration, but instead either defer creation of the object until after elaboration, or protect access to the object by setting a flag indicating if the object has been initialized or not. | Clariify the advice is only recomnmended if the language supports safe elaboration, otherwise either try to defer creation of the object until after elaboration (such via access using a pointer), or by putting guards around the object to ensure that the first access either initializes the object, or indicates failure if the object is not yet initialized. | Accept in principle. Replaced “elaboration” by “declaration”. |
| CA 010 |  | 06.23.01 |  | ed | "Each language provides rules of precedence and associativity, for each expression that operands bind to whichoperators. T" | I am unable to parse this sentence. Please reword. | Accept. Replace sentence by“Each language provides rules of precedence and associativity that determine for each expression which operands bind to which operators.” |
| CA 011 | 6th bullet | 06.25.05 |  | ed | The 6th bullet says avoid null statements, but I think the intent there (expressions that do nothing or are not executed, for instance) is getting confused with language defined null statements, which is actually recommended in 6.25.6, 1st bullet.  | Clarify that the 6th bullet is not referring to language defined null statements, but rather expressions or statements that do not have an effect, or something along those lines. | Accept. Change to “Avoid the use of statements that have no program effect.” |
| CA 012 |  | 06.25.06 |  | ed | "Languages should consider providing warnings for statements that are unlikely to be right such asstatements [without]{with} side effects. " | The logic of the sentence is backwards. “Without” should be replaced with “with”. | Accept.The sentence is confusing, and hence removed. |
| CA 013 |  | 06.26.03 |  | te | Another mechanism of failure, that could be mentioned is related to code bloat. A significant portion of dead code might effect the performance of the application (when loading the executable into memory, or perhaps cause linking problems by using up available code space on a tightly constrained embedded system. |  | Reject as not specific. The issue will be considered for a future revision. |
| CA 014 |  | 06.26.05 |  | ed | "When a developer identifies code that is dead because a conditional [consistently]{statically} evaluates " | Suggest replacing “consistently” with “statically” | Reject. The term “consistently” includes “statically”. Static evaluations can be detected by the compiler and dead code removed.. Non-static situations force the code to be left in the runtime system. |
| CA 015 |  | 06.34.04 |  | te | Probably variadic functions should be given special mention, such as printf(), where a format string is passed as an argument, and the logic of library, rather than the compiler, assumes that the number of parameters passed matches the format string. A recommendation for such languages that support variadic functions, could be to use altnernate calls, such as streaming, where the compiler can validate the correct parameters are passed, rather than relying on runtime functions. | Suggest mentioning variadic functions, and recommend avoiding their use if possible. | Variadic functions are not a subject of this vulnerability. It is covered under 6.64. |
| CA 016 |  | 06.39.04 |  | ed | "Languages {that} reclaim memory under programmer control can exhibit heap fragmentation and memoryleaks." | Add “that” | Accept. |
| CA 017 |  | 06.39.06 |  | ed | The word "pragma" is wrong font and size |  | Accept. |
| CA 018 |  | 06.40.03 |  | ed | " about the types it can legally be instantiated with." Avoid ending a sentence with a preposition. | Suggest replacing with => "about the parameterized types." | Accept. Used a different language construct. |
| CA 019 |  | 06.40.03 | Paragraph 3 | ed | There are two "For Example"s, which looks and reads as though one example was meant to replace the other. | Remove one of the examples, if only 1 is needed, or preface the second with something like, “As a second example,” | Accept. The second “for example” removed and the sentences connected. |
| CA 020 |  | 06.42.03 |  | ed | "of a class which [re]dispatches to the implementation" I dont think this applies only to redispatching, also to ordinary dispatching | Suggest generalizing and removing "re" | Accept. |
| CA 021 |  | 06.50.06 |  | te | "Provide a mechanism to determine which exceptions might be thrown by a called library routine." It is not clear this is good advice. Based on experience in other languages (in particular Java and C++), exception "signatures" have not generally provided the value expected for them. Here is an example paper describing some of the issues in the context of Java:   <http://literatejava.com/exceptions/checked-exceptions-javas-biggest-mistake/> <<http://literatejava.com/exceptions/checked-exceptions-javas-biggest-mistake/>>Here is a discussion of the problems with C++ exception specifications: <http://www.gotw.ca/publications/mill22.htm> <<http://www.gotw.ca/publications/mill22.htm>> | Consider whether this sentence should be removed or not. On the other hand, this is only advice, which the reader can ignore. There may be some who feel this is still good advice. Or perhaps clarification would be helpful to mention possible pitfalls of following this advice. | Accept. |
| PL 022 |  | Gereneral |  | ge | the document even though claiming to be language independent mostly depends on how Ada/C languages are implemented and mostly describes just them. Paper should either stay totally language independant and describe general concepts like floating point arithmetics problems etc. or only papers for specific languages should remain. | Leave only totally language independent topics in the document. | Not actionable. The authors of the comment Are invited to join the committee and provide advice on how to improve the document for future revisions. |

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