A proposal to add a utility class to represent expected monad (Revision 2)

Abstract

This paper is the 2nd revision of N4109 taking into account the feedback from the mailing lists and some fixes that were forgotten in revision 1.

Class template `expected<T,E>` proposed here is a type that may contain a value of type `T` or a value of type `E` in its storage space. `T` represents the expected value, `E` represents the reason explaining why it doesn't contain a value of type `T`. The interface and the rationale are based on `std::experimental::optional` and Haskell monads. We can consider `expected<T,E>` as a generalization of `optional<T>` providing in addition a monadic interface.

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History

Revision 2 - Revision of N4109 after discussion on the ML

- Fix default constructor to `T`. N4109 should change the default constructor to `T`, but there were some inconsistencies.
- Complete wording comparison.
- Adapted to last version of referenced proposals.
- Moved alternative designs from open questions to an Appendix.
- Moved already answered open points to a Rationale section.
- Moved open points that can be decided later to a future directions section.
- Complete wording hash.
- Add a section for adapting to `await`.
- Add a section in future work about a possible variadic.
- Fix minor typos.

Not done yet

- As `variant`, `expected` requires some properties in order to ensure the never-empty warranties. Add more on never-empty warranties and replace the wording.

Revision 1 - Revision of N4015 after Rapperswil feedback:

- Switch the expected class template parameter order from `expected<E,T>` to `expected<T,E>`.
- Make the unexpected value a salient attribute of the expected class concerning the relational operators.
- Removed open point about making expected and expected different classes.

Introduction

Class template `expected<T,E>` proposed here is a type that may contain a value of type `T` or a value of type `E` in its storage space. `T` represents the expected value, `E` represents the reason explaining why it doesn’t contain a value of type `T`, that is, the unexpected value. Its interface allows to query if the underlying value is either the expected value (of type `T`) or an unexpected value (of type `E`). The original idea comes from Andrei Alexandrescu C++ and Beyond 2012: Systematic Error Handling in C++
The interface and the rational are based on `std::experimental::optional` N3793 and Haskell monads. We can consider that `expected<T, E>` is a generalization of `optional<T>` providing in addition a monad interface and some specific functions associated to the unexpected type `E`. It requires no changes to core language, and breaks no existing code.

There is a related proposal for a class including an status and a optional value P0262R0, P0157R0 describes when to use each of the different error report mechanism.

## Motivation

Basicly, the two main error mechanisms are exceptions and return codes. Before further explanation, we should ask us what are the characteristics of a good error mechanism.

- **Error visibility**: Failure cases should appear throughout the code review. Because the debug can be painful if the errors are hidden.
- **Information on errors**: The errors should carry as most as possible information from their origin, causes and possibly the ways to resolve it.
- **Clean code**: The treatment of errors should be in a separate layer of code and as much invisible as possible. So the code reader could notice the presence of exceptional cases without stop his reading.
- **Non-Intrusive error**: The errors should not monopolize a communication channel dedicated to the normal code flow. They must be as discrete as possible. For instance, the return of a function is a channel that should not be exclusively reserved for errors.

The first and the third characteristic seem to be quite contradictory and deserve further explanation. The former points out that errors not handled should appear clearly in the code. The latter tells us that the error handling mustn't interfere with the code reading, meaning that it clearly shows the normal execution flow. A comparison between the exception and return codes is given in the next table.

<table>
<thead>
<tr>
<th>Visibility</th>
<th>Return error code</th>
</tr>
</thead>
<tbody>
<tr>
<td>Exception</td>
<td>Not visible without further analysis of the code. However, if an exception is thrown, we can follow the stack trace.</td>
</tr>
<tr>
<td>Return error code</td>
<td>Visible at the first sight by watching the prototype of the called function. However ignoring return code can lead to undefined results and it can be hard to figure out the problem.</td>
</tr>
<tr>
<td>Informations</td>
<td>Exceptions can be arbitrarily rich.</td>
</tr>
<tr>
<td>Clean code</td>
<td>Historically a simple integer. Nowadays, the header provides richer error code.</td>
</tr>
<tr>
<td>Clean code</td>
<td>Provides clean code, exceptions can be completely invisible for the caller.</td>
</tr>
<tr>
<td>Non-Intrusive</td>
<td>Force you to add, at least, a if statement after each function call.</td>
</tr>
<tr>
<td>Non-Intrusive</td>
<td>Proper communication channel.</td>
</tr>
<tr>
<td></td>
<td>Monopolization of the return channel.</td>
</tr>
</tbody>
</table>

### Expected class

We can do the same analysis for the `expected<T, E>` class and observe the advantages over the classic error reporting systems.

- **Error visibility**: It takes the best of the exception and error code. It's visible because the return type is `expected<T, E>` and the user cannot ignore the error case if he wants to retrieve the contained value.
- **Information**: Arbitrarily rich.
- **Clean code**: The monadic interface of expected provides a framework delegating the error handling to another layer of code. Note that `expected<T, E>` can also act as a bridge between an exception-oriented code and a nothrow world.
- **Non-Intrusive**: Use the return channel without monopolizing it.

It worthwhile mentioning the other characteristics of `expected<T, E>`:

- Associates errors with computational goals.
- Naturally allows multiple errors inflight.
- Teleportation possible.
- Across thread boundaries.
- Across no-throw subsystem boundaries.
- Across time: save now, throw later.
- Collect, group, combine errors.

### Use cases

#### Safe division

This example shows how to define a safe divide operation checking for divide-by-zero conditions. Using exceptions, we might write something like this:

```cpp
struct DivideByZero: public std::exception {...};

double safe_divide(double i, double j)
{
    if (j==0) throw DivideByZero();
    else return i / j;
}
```
With \texttt{expected\langle T, E \rangle}, we are not required to use exceptions, we can use \texttt{std::error\_condition} which is easier to introspect than \texttt{std::exception\_ptr} if we want to use the error. For the purpose of this example, we use the following enumeration (the boilerplate code concerning \texttt{std::error\_condition} is not shown):

\begin{verbatim}
enum class arithmetic_errc
{
    divide_by_zero, // 9/0 == ?
    not_integer_division // 5/2 == 2.5 (which is not an integer)
};
\end{verbatim}

Using \texttt{expected\langle double, error\_condition \rangle}, the code becomes:

\begin{verbatim}
expected\langle double, error\_condition \rangle safe_divide(double i, double j)
{
    if (j==0) return make\_unexpected(arithmetic\_errc::divide\_by\_zero); // (1)
    else return i / j; // (2)
}
\end{verbatim}

(1) The implicit conversion from \texttt{unexpected\_type\langle E \rangle} to \texttt{expected\langle T, E \rangle} and (2) from \texttt{T} to \texttt{expected\langle T, E \rangle} prevents using too much boilerplate code. The advantages are that we have a clean way to fail without using the exception machinery, and we can give precise information about why it failed as well. The liability is that this function is going to be tedious to use. For instance, the exception-based

\begin{verbatim}
function i + j/k is:
double f1(double i, double j, double k)
{
    return i + safe\_divide(j, k);
}
\end{verbatim}

but becomes using \texttt{expected\langle double, error\_condition \rangle}:

\begin{verbatim}
expected\langle double, error\_condition \rangle f1(double i, double j, double k)
{
    auto q = safe\_divide(j, k)
    if(q) return i + q;
    else return q;
}
\end{verbatim}

This example clearly doesn’t respect the “clean code” characteristic introduced above and the readability doesn’t differ much from the “C return code”. Hopefully, we can see \texttt{expected\langle T, E \rangle} through functional glasses as a monad. The code is cleaner using the member function \texttt{map}. This way, the error handling is not explicitly mentioned but we still know, thanks to the call to \texttt{map}, that something is going underneath and thus it is not as silent as exception.

\begin{verbatim}
expected\langle double, error\_condition \rangle f1(double i, double j, double k)
{
    return safe\_divide(j, k).map([k](double q) {
        return i + q;
    });
}
\end{verbatim}

The \texttt{map} member calls the continuation provided if expected contains a value, otherwise it forwards the error to the callee. Using lambda function might clutter the code, so here the same example using functor:

\begin{verbatim}
expected\langle double, error\_condition \rangle f1(double i, double j, double k)
{
    return safe\_divide(j, k).map(bind(plus, i, _1));
}
\end{verbatim}

We can use \texttt{expected\langle T, E \rangle} to represent different error conditions. For instance, with integer division, we might want to fail if the two numbers are not evenly divisible as well as checking for division by zero. We can overload our \texttt{safe\_divide} function accordingly:

\begin{verbatim}
expected\langle int, error\_condition \rangle safe\_divide(int i, int j)
{
    if (j==0) return make\_unexpected(arithmetic\_errc::divide\_by\_zero);
    if ((i | j) == 0) return make\_unexpected(arithmetic\_errc::not\_integer\_division);
    else return i / j;
}
\end{verbatim}

Now we have a division function for integers that possibly fail in two ways. We continue with the exception oriented
Now let's write this code using an `expected<T,E>` type and the functional `map` already used previously.

```
expected<int, error_condition> f(int i, int j, int k)
{
    return safe_divide(i, k).bind([-](int q1) {
        return safe_divide(j, k).map([-](int q2) {
            return q1 + q2;
        });
    });
}
```

The compiler will gently say he can convert an `expected<expected<int, error_condition>, error_condition>` to `expected<int, error_condition>` . This is because the member `map` wraps the result in `expected` and since we use twice the map member it wraps it twice. The `bind` member (do not confound with `std::bind`) wraps the result of the continuation only if it is not already wrapped. The correct version is as follow:

```
expected<int, error_condition> f(int i, int j, int k)
{
    return safe_divide(i, k).bind([-](int q1) {
        return safe_divide(j, k).bind([-](int q2) {
            return q1 + q2;
        });
    });
}
```

The error-handling code has completely disappeared but the lambda functions are a new source of noise, and this is even more important with `expected` variables. Propositions for a better monadic experience are discussed in section [Do-Notation], the subject is left open and is considered out of scope of this proposal.

**Error retrieval and correction**

The major advantage of `expected<T,E>` over `optional<T>` is the ability to transport an error, but we didn't come yet to an example that retrieve the error. First of all, we should wonder what a programmer do when a function call returns an error:

1. Ignore it.
2. Delegate the responsibility of error handling to higher layer.
3. Trying to resolve the error.

Because the first behavior might lead to buggy application, we won't consider it in a first time. The handling is dependent of the underlying error type, we consider the `exception_ptr` and the `error_condition` types.

We spoke about how to use the value contained in the `expected` but didn't discuss yet the error usage.

A first imperative way to use our error is to simply extract it from the `expected` using the `error()` member function. The following example shows a `divide2` function that return 0 if the error is `divide_by_zero`:

```
expected<int, error_condition> divide2(int i, int j)
{
    auto e = safe_divide(i, j);
    if (!e && e.error().value() == arithmetic_errc::divide_by_zero) {
        return 0;
    }
    return e;
}
```

This imperative way is not entirely satisfactory since it suffers from the same disadvantages than `value()`.

Again, a functional view leads to a better solution. The `catch_error` member calls the continuation passed as argument if the expected is erroneous.
An advantage of this version is to be coherent with the `bind` and `map` member functions. It also provides a more uniform way to analyze error and recover from some of these. Finally, it encourages the user to code its own “error-resolver” function and leads to a code with distinct treatment layers.

**Impact on the standard**

These changes are entirely based on library extensions and do not require any language features beyond what is available in C++ 14.

It requires however the `std::experimental::in_place_t` from N3793.

**Design rationale**

The same rationale described in N3672 for `optional<T>` applies to `expected<T,E>` and `expected<T, nullopt_t>` should behave almost as `optional<T>` with some exceptions. That is, we see `expected<T,E>` as `optional<T>` for which all the values of `E` collapse into a single value `nullopt`. In the following sections we present the specificities of the rationale in N3672 applied to `expected<T,E>`.

### Conceptual model of expected

`expected<T,E>` models a discriminated union of types `T` and `unexpected_type<E>`. `expected<T,E>` is viewed as a value of type `T` or value of type `unexpected_type<E>`, allocated in the same storage, along with the way of determining which of the two it is.

The interface in this model requires operations such as comparison to `T`, comparison to `E`, assignment and creation from either. It is easy to determine what the value of the expected object is in this model: the type it stores (`T` or `E`) and either the value of `T` or the value of `E`.

Additionally, within the affordable limits, we propose the view that `expected<T,E>` extends the set of the values of `T` by the values of type `E`. This is reflected in initialization, assignment, ordering, and equality comparison with both `T` and `E`. In the case of `optional<T>`, `T` cannot be a `nullopt_t`. As the types `T` and `E` could be the same in `expected<T,E>`, there is need to tag the values of `E` to avoid ambiguous expressions. The `make_unexpected(E)` function is proposed for this purpose. However `T` cannot be an `unexpected_type<E>` for a given `E`.

```cpp
expected<int, string> ei = 0;
expected<int, string> ej = 1;
expected<int, string> ek = make_unexpected(string());

ei = 1;
ej = make_unexpected(E());
ek = 0;

expected<int, string> ei = 0;
ex = make_unexpected(E());
ej = 0;
ex = 1;
```

### Initialization of `expected<T,E>`

In cases `T` and `E` have value semantic types capable of storing `n` and `m` distinct values respectively, `expected<T,E>` can be seen as an extended `T` capable of storing `n + m` values: these that `T` and `E` stores. Any valid initialization scheme must provide a way to put an expected object to any of these states. In addition, some `T`'s are not `CopyConstructible` and their expected variants still should be constructible with any set of arguments that work for `T`.

As in N3672, the model retained is to initialize either by providing an already constructed `T` or a tagged `E`. The default constructor required `T` to be default-constructible (as `expected<T>` should behave as `T` as much as possible).
There is still a minor issue as the default constructor of `std::exception_ptr` doesn’t contain an exception and so getting the value of a default constructed

```cpp
string s"STR";
expected_string.error_condition es[]; // requires Copyable<T>
expected_string.error_condition et -> s; // requires Copyable<T>
expected_string.error_condition ev = string"STR"; // requires Movable<T>
```

In order to create an unexpected object, the special function `make_unexpected` needs to be used:

```cpp
expected_string int ep(make_unexpected(-1)); // unexpected value, requires Movable<E>
expected_string int eq = make_unexpected(-1); // unexpected value, requires Movable<E>
```

As in N3672 and in order to avoid calling move/copy constructor of `T`, we use a “tagged” placement constructor:

```cpp
boost::MoveOnly error_condition eg; // expected value
boost::MoveOnly error_condition eh[]; // expected value
boost::MoveOnly error_condition el[in_place]; // calls MoveOnly{} in place
boost::MoveOnly error_condition ej[in_place, "arg"]; // calls MoveOnly("arg") in place
```

To avoid calling move/copy constructor of `E`, we use a “tagged” placement constructor:

```cpp
expected<int, string> ei(unexpect); // unexpected value, calls string{} in place
expected<int, string> ej(unexpect, "arg"); // unexpected value, calls string("arg") in place
```

An alternative name for `in_place` that is coherent with `unexpect` could be `expect`. Being compatible with `optional<T>` seems more important. So this proposal doesn’t propose such an `expect` tag.

Almost never-empty guaranty

**TODO:** revise this section as it cannot always be ensured

As `boost::variant<T,unexpected_type<E>>`, `expected<T,E>` ensures that it is never empty. All instances `v` of type `expected<T,E>` guarantee `v` has constructed content of one of the types `T` or `E`, even if an operation on `v` has previously failed.

This implies that expected may be viewed precisely as a union of exactly its bounded types. This “never-empty” property insulates the user from the possibility of undefined `expected` content and the significant additional complexity-of-use attendant with such a possibility.

In order to ensure this property the types `T` and `E` must satisfy some requirements as described in P0110R0. Given the nature of the parameter `E`, that is, to transport an error, it is expected that `is_nothrow_copy_constructible<E>` or at least `is_nothrow_move_constructible<E>`.

The default constructor

Similar data structure includes `optional<T>`, `variant<T1,...,Tn>` and `future<T>`. We can compare how they are default constructed.

- `std::optional<T>` default constructs to an optional with no value.
- `boost::variant<T1,...,Tn>` default constructs to the first type default constructible or it is ill-formed if none are default constructible.
- `std::future<T>` default constructs to an invalid future with no shared state associated, that is, no value and no exception.
- `std::optional<T>` default constructor is equivalent to `boost::variant<nullopt_t,T>`.

It raises several questions about `expected<T,E>`:

- Should the default constructor of `expected<T,E>` behave like `variant<T,unexpected_type<E>>` or as `variant<unexpected_type<T>,E>`?
- Should the default constructor of `expected<T,E>` behave like `optional<variant<T,E>>`?
- Should the default constructor of `expected<T,nullopt_t>` behave like `optional<T>`? If yes, how should behave the default constructor of `expected<T,E>`? As if initialized with `make_unexpected(E())`?
- Should `expected<T,E>` provide a default constructor at all? N3627 presents valid arguments against this approach, e.g. `array<expected<T,E>`, would not be possible.

Requiring `E` to be default constructible seems less constraining than requiring `T` to be default constructible (e.g. consider the `Date` example in N3527). With the same semantics `expected<Date,E>` would be `Regular` with a meaningful not-a-date state created by default.

There is still a minor issue as the default constructor of `std::exception_ptr` doesn’t contain an exception and so getting the value of a default constructed...
The authors consider the arguments in N3527 valid for `optional<T>`, however propose that `expected<T, E>` default constructor should behave as constructed with `T()` if `T` is default constructible.

### Conversion from `T`

An object of type `T` is implicitly convertible to an expected object of type `expected<T, E>`:

```cpp
expected<int, error_condition> ei = 1; // works
```

This convenience feature is not strictly necessary because you can achieve the same effect by using tagged forwarding constructor:

```cpp
expected<int, error_condition> ei{in_place, 1};
```

If the latter appears too cumbersome, one can always use function `make_expected` described below:

```cpp
expected<int> ei = make_expected(1);
auto ej = make_expected(1);
```

### Conversion from `E`

An object of type `E` is not convertible to an unexpected object of type `expected<T, E>` since `E` and `T` can be of the same type. The proposed interface uses a special tag `unexpect` and a special non-member `make_unexpected` function to indicate an unexpected state for `expected<T, E>`. It is used for construction and assignment. This might rise a couple of objections. First, this duplication is not strictly necessary because you can achieve the same effect by using the `unexpect` tag forwarding constructor:

```cpp
expected<string, int> exp1 = make_unexpected(1);
expected<string, int> exp2 = {unexpect, 1};
exp1 = make_unexpected(1);
exp2 = unexpect, 1;
```

While some situations would work with the `{unexpect, ...}` syntax, using `make_unexpected` makes the programmer's intention as clear and less cryptic. Compare these:

```cpp
expected<vector<int>, int> get1() {}
    return {unexpect, 1};

expected<vector<int>, int> get2() {
    return make_unexpected(1);

expected<vector<int>, int> get3() {
    return expected<vector<int>, int>{unexpect, 1};
}
```

The usage of `make_unexpected` is also a consequence of the adapted model for `expected`: a discriminated union of `T` and `unexpected_type<E>`. While `make_unexpected(E)` has been chosen because it clearly indicates that we are interested in creating an unexpected `expected<T, E>` (of unspecified type `T`), it could be also used to make a ready future with a specific error, but this is outside the scope of this proposal.

### Should we support the `exp2 = {}`?

Note also that the definition of the result type of `make_unexpected` has an explicitly deleted default constructor. This is in order to enable the reset idiom `exp2 = {}` which would otherwise not work due to the ambiguity when deducing the right-hand side argument.

*TODO: What is the meaning of exp2 = {}, now that expected defaults to T()?*

### Observers

In order to be as efficient as possible, this proposal includes observers with narrow and wide contracts. Thus, the `value()` function has a wide contract. If the expected object doesn't contain a value, an exception is thrown. However, when the user knows that the expected object is valid, the use of `operator*` would be more appropriate.

### Explicit conversion to bool

The rational described in N3672 for `optional<T>` applies to `expected<T, E>` and so, the following example combines initialization and value-checking in a boolean context.
if (expected<char, error_condition> ch = readNextChar()) {
    // ...
}

**has_value following P0032**

has_value has been added to follow [P0032R2].

**Accessing the contained value**

Even if `expected<T,E>` has not been used in practice for a while as Boost.Optional, we consider that following the same interface as `std::experimental::optional<T>` makes the C++ standard library more homogeneous.

The rational described in [N3672] for `optional<T>` applies to `expected<T,E>`.

**Dereference operator**

It was chosen to use indirection operator because, along with explicit conversion to bool, it is a very common pattern for accessing a value that might not be there:

```cpp
if (p) use(*p);
```

This pattern is used for all sort of pointers (smart or raw) and `optional`: it clearly indicates the fact that the value may be missing and that we return a reference rather than a value. The indirection operator has risen some objections because it may incorrectly imply `expected` and `optional` are a (possibly smart) pointer, and thus provides shallow copy and comparison semantics. All library components so far use indirection operator to return an object that is not part of the pointer's/iterator's value. In contrast, `expected` as well as `optional` indicts to the part of its own state. We do not consider it a problem in the design; it is more like an unprecedented usage of indirection operator. We believe that the cost of potential confusion is overweighted by the benefit of an intuitive interface for accessing the contained value.

We do not think that providing an implicit conversion to `T` would be a good choice. First, it would require different way of checking for the empty state; and second, such implicit conversion is not perfect and still requires other means of accessing the contained value if we want to call a member function on it.

Using the indirection operator for a object that doesn't contain a value is an undefined behavior. This behavior offers maximum runtime performance.

**Function value**

In addition to the indirection operator, we propose the member function `value` as in [N3672] that returns a reference to the contained value if one exists or throw an exception otherwise.

```cpp
void interact() {
    string s;
    cin >> s;
    expected<int, error> ei = str2int(s);
    try {
        process_int(ei.value());
    }
    catch (bad_expected_access<error>) {
        cout << "this was not a number."
    }
}
```

The exception thrown depends on the expected error type. By default it throws `bad_expected_access<E>` (derived from `std::logic_error`) which will contain the stored error. In the case of `expected<T, exception_ptr>`, it throws the exception stored in the `exception_ptr`. An approach enabling customization of this behavior is presented in the section [Customizing the exception thrown]: `bad_expected_access<E>` and `bad_optional_access` could inherit both from a `bad_access` exception derived from `logic_error`, but this is not proposed yet.

**Accessing the contained error**

Usually, accessing the contained error is done once we know the expected object has no value. This is why the `error()` function has a narrow contract: it works only if `!bool(*this)`.

```cpp
expected<int, errc> getIntOrZero(istream_range r) {
    auto r = getInt(); // won't throw
    if (!r && r.error() == errc::empty_stream) {
        return 0;
    }
    return r;
}
```
Conversion to the unexpected value

As the `error()` function, the `get_unexpected()` works only if the expected object has no value. It is used to propagate errors. Note that the following equivalences yield:

```cpp
f.get_unexpected() == make_unexpected(f.error());
f.get_unexpected() == expected<T, E>::unexpected; f.error();
```

This member is provided for convenience, it is further demonstrated in the next example:

```cpp
expected<pair<int, int>, error> getInRange(istream_range & r) {
    auto f = getInt(r);
    if (!f) return f.get_unexpected();
    auto m = matchedString("..", r);
    if (!m) return m.get_unexpected();
    auto l = getInt(r);
    if (!l) return l.get_unexpected();
    return std::make_pair(f, l);
}
```

`get_unexpected` is also provided for symmetry purpose. On one side, there is an implicit conversion from `expected_type<E>` to `expected<T, E>` and on the other side there is an explicit conversion from `expected<T, E>` to `expected_type<E>`. A more pleasant function manipulating error is `catch_error(F)` and is explained in the monadic operations section.

Function `value_or`

The function member `value_or()` has the same semantics than `optional` N3672 since the type of `E` doesn’t matter; hence we can consider that `E == nullopt_t` and the `optional` semantics yields.

Using the monadic error interface, we can achieve a similar behavior:

```cpp
auto x = getInt();
int x = *(x.value_or(0)); // identical to x.value_or(0);
```

Relational operators

As `optional`, one of the design goals of `expected` is that objects of type `expected<T, E>` should be valid elements in STL containers and usable with STL algorithms (at least if objects of type `T` and `E` are). Equally comparison is essential for `expected<T, E>` to model concept `Regular`. C++ does not have concepts yet, but being regular is still essential for the type to be effectively used with STL.

Ordering is essential if we want to store expected values in ordered associative containers. A number of ways of including the unexpected state in comparisons have been suggested. The ones proposed, have been crafted such that the axioms of equivalence and strict weak ordering are preserved: unexpected values stored in `expected<T, E>` are simply treated as additional values that are always different from `T`; these values are always compared as less than any value of `T` when stored in an `expected` object.

The main issue is how to compare the unexpected values between them. `operator==()` is defined for `exception_ptr`, using shallow semantics but there is no order between two `exception_ptr`.

```cpp
template<class T, class E>
constexpr bool operator==(const expected<T, E>& x, const expected<T, E>& y) {
    return x == y ? true : false;
}
template<class T, class E>
constexpr bool operator!=(const expected<T, E>& x, const expected<T, E>& y) {
    return x != y ? true : false;
}
```

If we follow the `optional<T>` semantics, two unexpected values should always be equal and do not compare. That is, `??` should be substituted by `false` and `??` by `true`. However considering all the unexpected value equals seems counterintuitive.

The alternative consists in forwarding the request to the respective `expected_type<E>` relational operators. That is, `??` should be substituted by `x.get_unexpected() < y.get_unexpected() and ??` by `x.get_unexpected() == y.get_unexpected()`.

But how to define the relational operators for `expected_type<E>`? We can forward the request to the respective `E` relational operators when `E` defines these operators
and follows the \texttt{optional<T>} semantics otherwise.

The case of \texttt{unexpected_type<std::exception_ptr>} could follow the \texttt{optional<T>} semantics as the shallow comparison is not very useful.

This limitation is one of the main motivations for having a user defined type with strict weak ordering. E.g. if the user know the exact types of the exceptions that can be thrown \texttt{E1, ... En}, the error parameter could be some kind of \texttt{variant<E1, ... En>} for which a strict weak ordering can be defined. If the user would like to take care of unknown exceptions something like \texttt{variant<monostate,E1, ... En>} would be a quite appropriated model.

\begin{verbatim}
expected<unsigned, int> e0[8];
expected<unsigned, int> e1[1];
expected<unsigned, int> eN[unexpect, -1];
assert (eN < e0);
assert (e0 < e1);
assert ([(eN < eN)])
assert ([(e1 < e1)])
assert (eN != e0);
assert (e0 != e1);
assert (eN == eN);
assert (e0 == e0);
\end{verbatim}

Unexpected values could have been as well considered greater than any value of \texttt{T}. The choice is a great degree arbitrary. We choose to stick to what \texttt{std::optional} does.

Given that both \texttt{unexpected_type<E>} and \texttt{T} are implicitly convertible to \texttt{expected<T,E>}, this implies the existence and semantics of mixed comparison between \texttt{expected<T,E>} and \texttt{T}, as well as between \texttt{expected} and \texttt{unexpected_type}:

\begin{verbatim}
assert (eN == make_unexpected(1));
assert (e0 != make_unexpected(1));
assert (eN != 1);
assert (e1 == 1);
assert (eN < 1);
assert (e0 > make_unexpected(1));
\end{verbatim}

Although it is difficult to imagine any practical use case of ordering relation between \texttt{expected<T,E>} and \texttt{unexpected_type<E>}, we still provide it for completeness sake.

The mixed relational operators, especially those representing order, between \texttt{expected<T,E>} and \texttt{T} have been accused of being dangerous. In code examples like the following, it may be unclear if the author did not really intend to compare two \texttt{T}’s.

\begin{verbatim}
auto count - get_expected_count();
if (count < 20) // or did you mean: *count < 20 ?
if (!count || *count < 20) // verbose, but unambiguous
\end{verbatim}

Given that \texttt{expected<T,E>} is comparable and implicitly constructible from \texttt{T}, the mixed comparison is there already. We would have to artificially create the mixed overloads only for them to cause controlled compilation errors. A consistent approach to prohibiting mixed relational operators would be to also prohibit the conversion from \texttt{T} or to also prohibit homogenous relational operators for \texttt{expected<T,E>} ; we do not want to do either, for other reasons discussed in this proposal. Also, mixed relational operations are available in \texttt{std::optional<T>} and we want to maintain the same behavior for \texttt{expected<T,nullopt_T>} and \texttt{optional<T>}. Mixed operators come as something natural when we consider the model “\texttt{T} with additional values”.

For completeness sake, we also provide ordering relations between \texttt{expected<T,E>} and \texttt{unexpected_type<E>}, even though we see no practical use case for them:

\begin{verbatim}
bool test(expected<unsigned, int> e)
{
    assert (e >= make_unexpected(1));
    assert ([(e < make_unexpected(1))]);
    assert (make_unexpected(1) <= e);
    return (e > make_unexpected(1));
}
\end{verbatim}

There exist two ways of implementing \texttt{operator>} for \texttt{expected} objects: use \texttt{T::operator>} or use \texttt{expected<T,E>::operator>}.

In case \texttt{T::operator>} and \texttt{T::operator<} are defined consistently, both above implementations are equivalent.

If the two operators are not consistent, the choice of implementation makes a difference. For relational operations, we choose to implement all in terms of \texttt{expected<T,E>::operator>} to be consistent with the choice taken for \texttt{std::optional}.

The same applies to the relational operators for \texttt{unexpected_type<E>}.

\section*{Modifiers}

\section*{Resetting the value}
Resetting the value of `expected<T,E>` is similar to `optional<T>` but instead of building a disengaged `optional<T>`, we build an erroneous `expected<T,E>`. Hence, the semantics and rationale is the same than in N3672.

**Tag `in_place`**

This proposal makes use of the “in-place” tag defined in N3793. This proposal provides the same kind of “in-place” constructor that forwards (perfectly) the arguments provided to `expected`’s constructor into the constructor of `T`. Hence, the semantics and rationale is the same than in N3672.

In order to trigger this constructor one has to use the tag `in_place`. We need the extra tag to disambiguate certain situations, like calling `expected`’s default constructor and requesting `T`’s default construction:

```cpp
expected<Big, error> eb{in_place, "1"}; // calls Big{"1"} in place (no moving)
expected<Big, error> ec{in_place}; // calls Big{} in place (no moving)
expected<Big, error> ed{}; // calls Big{} (expected state)
```

**Tag `unexpect`**

This proposal provides an “unexpect” constructor that forwards (perfectly) the arguments provided to `expected`’s constructor into the constructor of `E`. In order to trigger this constructor one has to use the tag `unexpect`.

We need the extra tag to disambiguate certain situations, notably if `T` and `E` are the same type.

```cpp
expected<Big, error> eb{unexpect, "1"}; // calls error{"1"} in place (no moving)
expected<Big, error> ec{unexpect}; // calls error{} in place (no moving)
```

In order to make the tag uniform an additional “expect” constructor could be provided but this proposal doesn’t propose it.

**Requirements on `T` and `E`**

Class template `expected` imposes little requirements on `T` and `E`: they have to be complete object type satisfying the requirements of `Destructible`. Each operations on `expected<T,E>` have different requirements and may be disable if `T` or `E` doesn’t respect these requirements. For example, `expected<T,E>`’s move constructor requires that `T` and `E` are `MoveConstructible`, `expected<T,E>`’s copy constructor requires that `T` and `E` are `CopyConstructible`, and so on. This is because `expected<T,E>` is a wrapper for `T` or `E` and it should resemble `T` as much as possible. If `T` is `EqualityComparable` then (and only then) we expect `expected<T,E>` to be `EqualityComparable`.

todo Refine the requirements, as this is not completely true, we need nothrow ...

**Expected references**

This proposal doesn’t include `expected` references as `optional` N3793 doesn’t include references neither.

**Expected void**

While it could seem weird to instantiate `optional` with `void`, it has more sense for `expected` as it conveys in addition, as `future<T>`, an error state.

**Making expected a literal type**

In N3672, they propose to make `optional` a literal type, the same reasoning can be applied to expected. Under some conditions, such that `T` and `E` are trivially destructible, and the same described for `optional`, we propose that `expected` be a literal type.

**Moved from state**

We follow the approach taken in `optional` N3672. Moving `expected<T,E>` do not modify the state of the source (valued or erroneous) of `expected` and the move semantics is up to `T` or `E`.

**IO operations**

For the same reasons than `optional` N3672 we do not add `operator<<` and `operator>>` IO operations.

**Monadic operations**

A monadic interface is not optional if we don’t want to fall back in the problems of the old “C return code”. The example section shows how these operations are important to `expected`. The member function `map` and `bind` find their roots in the category theory if we consider `expected` as a functor and a monad.
Functor map

The operation `map` consider `expected` as a Functor and just apply a function on the contained value, if any. The types of the two overloads are presented using a functional notation and the `[]` represents a context in which the value `T` or `U` is contained. The current context is `expected<T,E>` and thus `[T]` is equivalent to `expected<T,E>` for a specific `E`.

\[(T \to U) \to [U]\]

Whatever the return type of the continuation, we observe that it is always wrapped into a context. The monadic bind do it differently.

Monadic bind

A Monad is defined with a type constructor and two operations `make` and `bind`. The type constructor simply builds a monad for a specific type, in the C++ jargon it is referred to template instantiation (we build `expected` from a type Value and Error).

The `make` operation wraps a value of type `T` inside a context `[T]`. In C++ we can consider the constructors as a `make` operation.

Finally, the `bind` operation is similar to `map` but doesn’t wrap the value if the function already wraps it up.

The functional signature of bind can be described as follow:

\[(T \to [U]) \to [U]\]

If a do-notation is introduced in C++, as proposed in section [Do-Notation], these operations can become a powerful abstraction, they have been proven very useful in Haskell. For example, a similar interface could be used with `optional`.

then function

The last operation has no direct counterpart in functional language and is inspired from N3857 proposing some improvements to `std::future<T>`. The functional signature is as follow:

\[[T] \to [U] \to [U] \to [U]\]

It has a conditional wrapping strategy: it doesn’t wrap if the continuation already wraps it up.

Exception thrown in the continuation

Currently, the exceptions thrown in the continuations are not caught.

catch_error function

This last member function is used when we want to use or recover from an error. When chaining multiple bind or map operations we don’t know if the operations have succeeded. A common way is thus to add a catch_error at the end and act in consequence.

```
getInt().map([](int i){return i * 2;})
 .map(integer_divide_by_2)
 .catch_error(log_error);
```

Here the last operation is simply used to log the error but the `catch_error` also accepts function that try to recover from a previous error.

```
getInt().map([](int i){return i * 2;})
 .map(integer_divide_by_2)
 .catch_error([](auto e) return 0;);
```

This last example shows we can return a new value from the continuation passed to catch_error. The `catch_error` member doesn’t catch exceptions that could be thrown by the continuation. Since we already try to recover from an error it makes little sense to prevent the user to launch an exception.

unwrap function

In some scenarios, you might want to create an `expected` that returns another `expected`, resulting in nested `expected`. It is possible to write simple code to unwrap the outer expected and retrieve the nested expected and its result with the current interface as in:
template <class T, class E>
expected<T,E> unwrap(expected::expected<T,E,E> ee) {
    if (ee) return *ee;
    return ee.get_unexpected();
}

// class expected
template <class T, class E>
expected<T,E> unwrap(expected::expected<T,E>E e) {
    return e;
}

We could add such a function to the standard, either as a free function or as a member function.

# Open points

**Should** `expected<T, exception_ptr>` **be equality comparable?**

This proposal makes `expected<T, exception_ptr>` equality comparable making all the unexpected values equals as `exception_ptr` equality comparison is shallow and doesn't provide relational operators.

Should `expected<T, exception_ptr>` not be equality comparable?

Should `expected<T, exception_ptr>` be equality comparable using shallow comparison?

**Should** `expected<T,E>` **throw E instead of bad_expected_access<E>**?

As any type can be thrown as an exception, should `expected<T,E>` throw `E` instead of `bad_expected_access<E>`?

If yes, should `optional<T>` throw `nullopt_t` instead of `bad_optional_access` to be coherent?

Should `expected<T, exception_ptr>` throw `exception_ptr` instead of the stored exception to be coherent?

**Should** `expected<T,E>` **be convertible from E when E is not convertible to T**?

The implicit conversion from `E` has been forbidden to avoid ambiguity when `E` and `T` are the same type. However when `E` and `T` are not convertible there wouldn't any ambiguity.

Should the implicit conversion be allowed in this case?

**Should** map/bind/then/catch_error **catch the exceptions throw by the continuation?**

It is easy to catch the exceptions when the type is `expected<T>`. However, doing it for `expected<T,E>` needs a conversion from the current exception and the error `E`.

This proposal requires that the continuation doesn’t throw exceptions as we don’t have a general solution.

Should `expected<T>::map/bind/then/catch_error` catch the exceptions and propagate them on the result?

Should `expected<T,E>::map/bind/then/catch_error` catch the exceptions and propagate them on the result by doing a transformation from the exception to the error?

If yes, how to configure it?

**Do we need a `expected<T,E>::error_or` function?**

It has been argued that the error should be always available and that often there is a success value associated to the error.

`expected<T,E>` would be seen more like something like `struct {E, optional<T>}`.

The following code was show as use case

```cpp
auto e = function();
switch (e.status())
    success: ..., break;
    too_green: ..., break;
    too_pink: ..., break;
```

With the current interface the user could be tempted to do
auto e = function();
if (e)
  /*success:*/ ....;
else
  switch (e.error())
    case too_green: ....; break;
    case too_pink: ....; break;

This could be done with the current interface as follows

auto e = function();
switch (error_or(e, success))
  success: ....; break;
  too_green: ....; break;
  too_pink: ....; break;

where

    template <class E, class T>
    E error_or(expected<T,E> const&, E err) {
      if(e) return err;
      else return error();
    }

Do we need to add such an `error_or` function as member?

Do we need a `expected<T,E>::has_error` function?

An other use case which could look much uglier is if the user had to test for whether or not there was a specific error code.

    e = function();
    while (e.status == timeout) {
      sleep(delay);
      delay *= 2;
      e = function();
    }

Here we have a value or a hard error. This use case would need to use something like `has_error`

    e = function();
    while (has_error(e, timeout) )
    {
      sleep(delay);
      delay *= 2;
      e = function();
    }

where

    template <class T, class E>
    bool has_error(expected<T,E> const&, E err) {
      if (e) return false;
      else return error() == err;
    }

Do we want to add such a `has_error` function as member?

Proposed Wording

The proposed changes are expressed as edits to N4564 the Working Draft - C++ Extensions for Library Fundamentals V2. The wording has been adapted from the section "Optional objects".

General utilities library

---------------------------------------------------------------------------- Insert a new section.----------------------------------------------------------------------------
X.Y Unexpected objects [unexpected]

X.Y.1 In general [unexpected.general]

This subclause describes class template `unexpected_type` that wraps objects intended as unexpected. This wrapped unexpected object is used to be implicitly convertible to other objects.

X.Y.2 Header synopsis [unexpected.synop]

```cpp
namespace std {
namespace experimental {
inline namespace fundamentals_v3 {
    // X.Y.3, Unexpected object type
    template <class E>
    class unexpected_type;
    // X.Y.4, Unexpected exception_ptr specialization
    template <>
    class unexpected_type::exception_ptr specialization
    template <>
    class unexpected_type::exception_ptr;
    // X.Y.5, Unexpected factories
    template <class E>
    constexpr unexpected_type decay_t<T=E> make_unexpected(E&& v);
    unexpected_type::std::exception_ptr make_unexpected_from_current_exception();
    // X.Y.6, unexpected_type relational operators
    template <class E>
    constexpr bool operator==(const unexpected_type::E&, const unexpected_type::E&);
    template <class E>
    constexpr bool operator!=(const unexpected_type::E&, const unexpected_type::E&);
    template <class E>
    constexpr bool operator>(const unexpected_type::E&, const unexpected_type::E&);
    template <class E>
    constexpr bool operator<(const unexpected_type::E&, const unexpected_type::E&);
    template <class E>
    constexpr bool operator&>(const unexpected_type::E&, const unexpected_type::E&);
    template <class E>
    constexpr bool operator&(const unexpected_type::E&, const unexpected_type::E&);
    // X.Y.7, unexpected_type relational operators
    constexpr bool operator==(const unexpected_type::exception_ptr&, const unexpected_type::exception_ptr&);
    constexpr bool operator!=(const unexpected_type::exception_ptr&, const unexpected_type::exception_ptr&);
    constexpr bool operator>(const unexpected_type::exception_ptr&, const unexpected_type::exception_ptr&);
    constexpr bool operator<(const unexpected_type::exception_ptr&, const unexpected_type::exception_ptr&);
    constexpr bool operator&(const unexpected_type::exception_ptr&, const unexpected_type::exception_ptr&);
    constexpr bool operator&>(const unexpected_type::exception_ptr&, const unexpected_type::exception_ptr&);
}}
```

A program that needs the instantiation of template `unexpected_type` for a reference type or `void` is ill-formed.

X.Y.3 Unexpected object type [unexpected.object]

```cpp
template <class E=std::exception_ptr>
class unexpected_type {
public:
    unexpected_type() = delete;
    constexpr explicit unexpected_type(E const&);
    constexpr explicit unexpected_type(E&&);
    constexpr E& const& value() const;
    constexpr E& & value();
private:
    E val; // exposition only
};
```
constexpr explicit unexpected_type(E const &);

Effects: Build an unexpected by copying the parameter to the internal storage val.

constexpr explicit unexpected_type(E & &);

Effects: Build an unexpected by moving the parameter to the internal storage val.

constexpr E const & value();
constexpr E const & value() const;

Returns: val.

X.Y.4 Unexpected exception ptr specialization [unexpected.exceptionptr]

template <>
class unexpected_type::exception_ptr { 
public:
    unexpected_type() = delete;
    explicit unexpected_type(exception_ptr const &);
    explicit unexpected_type(exception_ptr & &);
    template <class E>
    explicit unexpected_type(E);
    exception_ptr & value();
    exception_ptr const & value() const;
private:
    exception_ptr val; // exposition only
};

constexpr explicit unexpected_type(exception_ptr const &);

Effects: Build an unexpected by copying the parameter to the internal storage val.

constexpr explicit unexpected_type(exception_ptr & &);

Effects: Build an unexpected by moving the parameter to the internal storage val.

constexpr explicit unexpected_type(E e);

Effects: Build an unexpected storing the result of val(make_exception_ptr(e)). constexpr exception_ptr const & value() const;

exception_ptr & value();
exception_ptr const & value() const;

Returns: val.

X.Y.5 Factories [unexpected.factories]

template <class E>
constexpr unexpected_type::decay_t<E> make_unexpected(E & & v);

Returns: unexpected_type<decay_t<E>>(v).

constexpr unexpected_type::std::exception_ptr make_unexpected_from_current_exception();

Returns: unexpected_type<std::exception_ptr>(std::current_exception()).

X.Y.8 Unexpected Relational operators [unexpected.type.relationalop]

template <class E>
constexpr bool operator==(const unexpected_type<E> & x, const unexpected_type<E> & y);

Requires: E shall meet the requirements of EqualityComparable.

Returns: x.value() != y.value().
Remarks: Specializations of this function template, for which \( x.value() \neq y.value() \) is a core constant expression, shall be constexpr functions.

```cpp
template <class E>
constexpr bool operator!=((const unexpected_type<E>& x, const unexpected_type<E>& y));
```

Returns: \((x == y)\).

```cpp
template <class E>
constexpr bool operator<(const unexpected_type<E>& x, const unexpected_type<E>& y);
```

Requires: \(x.value() < y.value()\).

Returns: \(x.value() < y.value()\).

Remarks: Specializations of this function template, for which \( x.value() < y.value() \) is a core constant expression, shall be constexpr functions.

```cpp
template <class E>
constexpr bool operator>(const unexpected_type<E>& x, const unexpected_type<E>& y);
```

Returns: \(y < x\).

```cpp
template <class E>
constexpr bool operator<=(const unexpected_type<E>& x, const unexpected_type<E>& y);
```

Returns: \!(y < x)\).

```cpp
template <class E>
constexpr bool operator>=(const unexpected_type<E>& x, const unexpected_type<E>& y);
```

Returns: \!(x < y)\).

X.Y.7 Unexpected Relational operators [unexpected_type.relationlop]

```cpp
constexpr bool operator==(const unexpected_type_exception_ptr&E x, const unexpected_type_exception_ptr&E y);
```

Returns: \(x.value() \neq y.value()\).

Remarks: Specializations of this function template, for which \( x.value() \neq y.value() \) is a core constant expression, shall be constexpr functions.

```cpp
constexpr bool operator!=(const unexpected_type_exception_ptr&E x, const unexpected_type_exception_ptr&E y);
```

Returns: \((x == y)\).

```cpp
constexpr bool operator<(const unexpected_type_exception_ptr&E x, const unexpected_type_exception_ptr&E y);
```

Returns: \(false\).

```cpp
constexpr bool operator>(const unexpected_type_exception_ptr&E x, const unexpected_type_exception_ptr&E y);
```

Returns: \(false\).

```cpp
constexpr bool operator<=(const unexpected_type_exception_ptr&E x, const unexpected_type_exception_ptr&E y);
```

Returns: \(x == y\).

```cpp
constexpr bool operator>=(const unexpected_type_exception_ptr&E x, const unexpected_type_exception_ptr&E y);
```

Returns: \(x == y\).
X.Z Expected objects [expected]

X.Z.1 In general [expected.general]

This sub-clause describes class template expected that represents expected objects. An expected object is an object that contains the storage for another object and manages the lifetime of this contained object. Alternatively, it could contain the storage for another unexpected object. The contained object may not be initialized after the expected object has been initialized, and may not be destroyed before the expected object has been destroyed. The initialization state of the contained object is tracked by the expected object.

X.Z.2 Header <experimental/expected> synopsis [expected.synop]

```cpp
namespace std {
namespace experimental {
inline namespace fundamentals_v3 {
// X.Z.4, expected for object types
    template <class T, class E> exception_ptr
        class expected {

// X.Z.5, Specialization for void.
    template <class E>
        class expected<void, E> {

// X.Z.6, unexpected tag
    struct unexpected_t {
        unexpected_t() = delete;
    };

constexpr unexpected_t unexpected("implementation defined");

// X.Z.7, class bad_expected_access
    class bad_expected_access;

// X.Z.8, Expected relational operators
    template <class T, class E>
       constexpr bool operator==(const expected<T,E>&, const expected<T,E>&);
    template <class T, class E>
       constexpr bool operator!=(const expected<T,E>&, const expected<T,E>&);
    template <class T, class E>
       constexpr bool operator<(const expected<T,E>&, const expected<T,E>&);
    template <class T, class E>
       constexpr bool operator>(const expected<T,E>&, const expected<T,E>&);
    template <class T, class E>
       constexpr bool operator<=(const expected<T,E>&, const expected<T,E>&);
    template <class T, class E>
       constexpr bool operator>=(const expected<T,E>&, const expected<T,E>&);

// X.Z.9, Comparison with T
    template <class T, class E>
        constexpr bool operator==(const expected<T,E>&, const T&);
    template <class T, class E>
        constexpr bool operator!=(const expected<T,E>&, const T&);
    template <class T, class E>
        constexpr bool operator<(const expected<T,E>&, const T&);
    template <class T, class E>
        constexpr bool operator>(const expected<T,E>&, const T&);
    template <class T, class E>
        constexpr bool operator<=(const expected<T,E>&, const T&);
    template <class T, class E>
        constexpr bool operator>=(const expected<T,E>&, const T&);

// X.Z.10, Comparison with unexpected type E
    template <class T, class E>
        constexpr bool operator==(const expected<T,E>&, const unexpected_type<E>&);
    template <class T, class E>
        constexpr bool operator!=(const expected<T,E>&, const unexpected_type<E>&);
    template <class T, class E>
        constexpr bool operator<(const expected<T,E>&, const unexpected_type<E>&);
    template <class T, class E>
        constexpr bool operator>(const expected<T,E>&, const unexpected_type<E>&);
    template <class T, class E>
        constexpr bool operator<=(const expected<T,E>&, const unexpected_type<E>&);
    template <class T, class E>
        constexpr bool operator>=(const expected<T,E>&, const unexpected_type<E>&);
}
}
```

A program that necessitates the instantiation of template `expected<T,E>` with `T` for a reference type or for possibly cv-qualified types `in_place_t`, `unexpect_t` or `unexpected_type<E>` is ill-formed.

X.Z.3 Definitions [expected.defs]

An instance of `expected<T,E>` is said to be valued if it contains a value of type `T`. An instance of `expected<T,E>` is said to be unexpected if it contains an object of type `E`.

X.Y.4 expected for object types [expected.object]
template <class U, class... Args>
  constexpr explicit expected(in_place_t, initializer_list<U>., Args&&...);
constexpr expected(unexpected_type_E const&);

// X.Z.4.2, destructor
constexpr expected();

// X.Z.4.3, assignment
expected& operator=(const expected&);
expected& operator=(expected&&) noexcept(see below);

// X.Z.4.4, swap
void swap(expected&); swap(expected&&) noexcept(see below);

// X.Z.4.5, observers
constexpr T const* operator->() const;
T* operator->()&&;
constexpr T const& operator*() const&;
T& operator*()&;
constexpr T&& operator*()&&;
constexpr explicit operator bool() const noexcept;
constexpr T T const& value() const&;
T& value()&&;
constexpr T&& value()&&;
constexpr T const& error() const&;
E& error()&;
constexpr E&& error()&&;
constexpr unexpected_type_E get_unexpected() const;

// X.Z.4.6, factories
template <typename Ex, typename F>
  expected<T,E> catch_exception(F&& f);

// exposition only
private:
  bool has_value; // exposition only
  union {
    value_type val; // exposition only
    error_type err; // exposition only
  };
T and E shall be an object type and shall satisfy the requirements of Destructible.

X.Z.4.1 Constructors [expected.object.ctor]

### constexpr expected() noexcept(“see below”);

**Effects:** Initializes the contained value as if direct-non-list-initializing an object of type T with the expression T().

**Postconditions:** bool(*this).

**Throws:** Any exception thrown by the default constructor of T.

**Remarks:** The expression inside noexcept is equivalent to: is_nothrow_default_constructible<T>::value.

**Remarks:** This signature shall not participate in overload resolution unless is_default_constructible<T>::value.

### expected(const expected& rhs);

**Effects:** If bool(rhs) initializes the contained value as if direct-non-list-initializing an object of type T with the expression *rhs.

If !bool(rhs) initializes the contained value as if direct-non-list-initializing an object of type E with the expression rhs.error().

**Postconditions:** bool(rhs) == bool(*this).

**Throws:** Any exception thrown by the selected constructor of T or E.

**Remarks:** This signature shall not participate in overload resolution unless is_copy_constructible<T>::value and is_copy_constructible<E>::value.

### expected(expected && rhs) noexcept(“see below”);

**Effects:** If bool(rhs) initializes the contained value as if direct-non-list-initializing an object of type T with the expression std::move(*rhs).

If !bool(rhs) initializes the contained value as if direct-non-list-initializing an object of type E with the expression std::move(rhs.error()).

**Postconditions:** bool(rhs) == bool(*this) and bool(rhs) is unchanged.

**Throws:** Any exception thrown by the selected constructor of T or E.

**Remarks:** The expression inside noexcept is equivalent to: is_nothrow_move_constructible<T>::value == type and is_nothrow_move_constructible<E>::value.

**Remarks:** This signature shall not participate in overload resolution unless is_move_constructible<T>::value and is_move_constructible<E>::value.

### constexpr expected(const T& v);

**Effects:** Initializes the contained value as if direct-non-list-initializing an object of type T with the expression v.

**Postconditions:** bool(*this).

**Throws:** Any exception thrown by the selected constructor of T.

**Remarks:** If T’s selected constructor is a constexpr constructor, this constructor shall be a constexpr constructor.

**Remarks:** This signature shall not participate in overload resolution unless is_copy_constructible<T>::value.

### constexpr expected(T&& v);

**Effects:** Initializes the contained value as if direct-non-list-initializing an object of type T with the expression std::move(v).

**Postconditions:** bool(*this).

**Throws:** Any exception thrown by the selected constructor of T.

**Remarks:** If T’s selected constructor is a constexpr constructor, this constructor shall be a constexpr constructor.

**Remarks:** This signature shall not participate in overload resolution unless is_move_constructible<T>::value.

### template <class... Args>

castexpr explicit expected(in_place_t, Args&&... args);

**Effects:**

**Postconditions:**

**Throws:**

**Remarks:**

**Remarks:**

**Remarks:**
Effects: Initializes the contained value as if direct-non-list-initializing an object of type `T` with the arguments `std::forward<Args>(args)...`.

Postconditions: `bool(*this)`.

Throws: Any exception thrown by the selected constructor of `T`.

Remarks: If `T`’s constructor selected for the initialization is a constexpr constructor, this constructor shall be a constexpr constructor.

Remarks: This signature shall not participate in overload resolution unless `is_constructible<T, Args&&...>::value`.

```cpp
template<class U, class... Args>
constexpr explicit expected(in_place_t, initializer_list<U> il, Args&&... args);
```

Effects: Initializes the contained value as if direct-non-list-initializing an object of type `T` with the arguments `il, std::forward<Args>(args)...`.

Postconditions: `bool(*this)`.

Throws: Any exception thrown by the selected constructor of `T`.

Remarks: The function shall not participate in overload resolution unless: `is_constructible<T, initializer_list<U>&, Args&&...>::value`.

If `T`’s constructor selected for the initialization is a constexpr constructor, this constructor shall be a constexpr constructor.

Remarks: This signature shall not participate in overload resolution unless `is_constructible<T, initializer_list<U>&, Args&&...>::value`.

```cpp
constexpr expected(unexpected_type<E> const& e);
```

Effects: Initializes the unexpected value as if direct-non-list-initializing an object of type `E` with the expression `e.value()`.

Postconditions: `! bool(*this)`.

Throws: Any exception thrown by the selected constructor of `E`.

Remark: If `E`’s selected constructor is a constexpr constructor, this constructor shall be a constexpr constructor.

Remark: This signature shall not participate in overload resolution unless `is_copy_constructible<E>::value`.

```cpp
constexpr expected(unexpected_type<E>&& e);
```

Effects: Initializes the unexpected value as if direct-non-list-initializing an object of type `E` with the expression `std::move(e.value())`.

Postconditions: `! bool(*this)`.

Throws: Any exception thrown by the selected constructor of `E`.

Remark: If `E`’s selected constructor is a constexpr constructor, this constructor shall be a constexpr constructor.

Remark: This signature shall not participate in overload resolution unless `is_move_constructible<E>::value`.

```cpp
template<class... Args>
constexpr explicit expected(unexpect_t, Args&&... args);
```

Effects: Initializes the unexpected value as if direct-non-list-initializing an object of type `E` with the arguments `std::forward<Args>(args)...`.

Postconditions: `! bool(*this)`.

Throws: Any exception thrown by the selected constructor of `E`.

Remarks: If `E`’s constructor selected for the initialization is a constexpr constructor, this constructor shall be a constexpr constructor.

Remarks: This signature shall not participate in overload resolution unless `is_constructible<E, Args&&...>::value`.

```cpp
template<class U, class... Args>
constexpr explicit expected(unexpect_t, initializer_list<U> il, Args&&... args);
```

Effects: Initializes the unexpected value as if direct-non-list-initializing an object of type `E` with the arguments `il, std::forward<Args>(args)...`.

Postconditions: `! bool(*this)`.

Throws: Any exception thrown by the selected constructor of `E`.

Remarks: The function shall not participate in overload resolution unless: `is_constructible<E, initializer_list<U>&, Args&&...>::value`.
If E’s constructor selected for the initialization is a constexpr constructor, this constructor shall be a constexpr constructor.

Remarks: This signature shall not participate in overload resolution unless is_constructible<E, initializer_list<U>&, Args&&...>::value.

X.2.4.2 Destructor [expected.object.dtor]

```cpp
 EXPECTED();
```

Effects: If is_trivially_destructible<T>::value != true and bool(*this), calls val->T::~T(). If is_trivially_destructible<E>::value != true and ! (bool(*this), calls err->E::~E()).

Remarks: If is_trivially_destructible<T>::value and is_trivially_destructible<E>::value then this destructor shall be a trivial destructor.

X.2.4.3 Assignment [expected.object.assign]

```cpp
expected<T,E&> operator=(const expected<T,E&> rhs);
```

Requires: is_nothrow_copy_constructible<E>::value && is_nothrow_move_constructible<E>::value is true

Effects:

- If bool(*this) and bool(rhs),
  - assigns *rhs to the contained value val, otherwise
  - If bool(*this) and ! bool(rhs),
    - assigns rhs.error() to the contained value err, otherwise
- If bool(*this) and ! bool(rhs),
  - destroys the contained value by calling val->T::~T(),
    - initializes the contained value as if direct-non-list-initializing an object of type T with rhs.error(), otherwise ! (bool(*this) and bool(rhs))
- If is_nothrow_copy_constructible<T>::value
  - destroys the contained value by calling err->E::~E(),
    - initializes the contained value as if direct-non-list-initializing an object of type T with *rhs, otherwise
- If is_nothrow_move_constructible<T>::value
  - constructs a new T tmp on the stack from *rhs,
    - destroys the contained value by calling err->E::~E(),
    - initializes the contained value as if direct-non-list-initializing an object of type T with tmp, otherwise
- If is_nothrow_move_constructible<E>::value
  - move constructs a new E tmp on the stack from this.error() (which can’t throw as A is nothrow-move-constructible),
    - destroys the contained value by calling err->E::~E(),
    - initializes the contained value as if direct-non-list-initializing an object of type T with *rhs. Either,
      - The constructor didn’t throw, so mark the expected as holding a T (which can't throw), or
      - The constructor did throw, so move-construct the E from the stack tmp back into the expected storage (which can't throw as E is nothrow-move-constructible), and rethrow the exception.

Returns: *this.

Postconditions: bool(rhs) == bool(*this).

Exception Safety: If any exception is thrown, the values of bool(*this) and bool(rhs) remain unchanged.

If an exception is thrown during the call to T’s copy constructor, no effect.

If an exception is thrown during the call to T’s copy assignment, the state of its contained value is as defined by the exception safety guarantee of T’s copy assignment.

If an exception is thrown during the call to E’s copy constructor, no effect.

If an exception is thrown during the call to E’s copy assignment, the state of its contained value is as defined by the exception safety guarantee of E’s copy assignment.

Remarks: TODO reword this part This signature shall not participate in overload resolution unless is_copy_constructible<T>::value and is_copyAssignable<T>::value and is_copy_constructible<E>::value and is_copy_assignable<E>::value.

```cpp
expected<T,E&> operator=(expected<T,E&> rhs) noexcept(/" see below");
```
TODO: reword this operation

Effects: If `bool(*this)` and `rhs` is values, assigns `std::move(*rhs)` to the contained value `val`, otherwise if `bool(*this)` and `bool(rhs)`, destroys the contained value by calling `val->T::~T()` and initializes the contained value as if direct-non-list-initializing an object of type `T` with `rhs.error()`; otherwise if `! bool(*this)` and `bool(rhs)`, assigns `std::move(rhs.error())` to the contained value `err`, otherwise if `! bool(*this)` and `bool(rhs)`, destroys the contained value by calling `err->E::~E()` and initializes the contained value as if direct-non-list-initializing an object of type `E` with `rhs.error()`.

Returns: `*this`.

Postconditions: `bool(rhs) == bool(*this)`.

Remarks: The expression inside noexcept is equivalent to:

- `is_nothrow_move_assignable<T>::value && is_nothrow_move_constructible<T>::value && is_nothrow_move_assignable<E>::value && is_nothrow_move_constructible<E>::value && is_nothrow_move_assignable<T>::value && is_nothrow_move_constructible<T>::value && is_nothrow_move_assignable<E>::value && is_nothrow_move_constructible<E>::value && is_nothrow_move_assignable<T>::value && is_nothrow_move_constructible<T>::value && is_nothrow_move_assignable<E>::value && is_nothrow_move_constructible<E>::value`.

Exception Safety: If any exception is thrown, the values of `bool(*this)` and `bool(rhs)` remain unchanged. If an exception is thrown during the call to `T`'s move constructor, the state of `rhs.val` is determined by exception safety guarantee of `T`'s move constructor. If an exception is thrown during the call to `T`'s move assignment, the state of `val` and `rhs.val` is determined by exception safety guarantee of `T`'s move assignment. If an exception is thrown during the call to `E`'s move constructor, the state of `rhs.err` is determined by exception safety guarantee of `E`'s move constructor. If an exception is thrown during the call to `E`'s move assignment, the state of `err` and `rhs.err` is determined by exception safety guarantee of `E`'s move assignment.

Remarks: This signature shall not participate in overload resolution unless `is_move_constructible<T>::value and is_move_assignable<T>::value and is_move_constructible<E>::value and is_move_assignable<E>::value`.

```cpp
template <class U
    expected T E & operator=(U & v);
```

TODO: reword this operation

Effects: If `bool(*this)` assigns `std::forward<U>(v)` to the contained value; otherwise destroys the contained value by calling `err->E::~E()` and initializes the unexpected value as if direct-non-list-initializing an object of type `T` with `std::forward<U>(v)`.

Returns: `*this`.

Postconditions: `bool(*this)`.

Exception Safety: If any exception is thrown, `bool(*this)` remains unchanged. If an exception is thrown during the call to `E`'s constructor, the state of `v` is determined by exception safety guarantee of `E`'s constructor. If an exception is thrown during the call to `E`'s assignment, the state of `err` and `e` is determined by exception safety guarantee of `E`'s assignment.

Remarks: This signature shall not participate in overload resolution unless `is_constructible<T, U>::value and is_assignable<T &, U>::value`.

[Note: The reason to provide such generic assignment and then constraining it so that effectively `T == U` is to guarantee that assignment of the form `o = {}` is unambiguous. --end note]

```cpp
expected T E & operator=(unexpected_type<E &> e);
```

TODO: reword this operation

Effects: If `! bool(*this)` assigns `std::forward<E>(e.value())` to the contained value; otherwise destroys the contained value by calling `val->T::~T()` and initializes the contained value as if direct-non-list-initializing object of type `T` with `std::forward<unexpected_type<E>>(e).value()`.

Returns: `*this`.

Postconditions: `! bool(*this)`.

Exception Safety: If any exception is thrown, value of `val` remains unchanged. If an exception is thrown during the call to `T`'s constructor, the state of `v` is determined by exception safety guarantee of `T`'s constructor. If an exception is thrown during the call to `T`'s assignment, the state of `val` and `v` is determined by exception safety guarantee of `T`'s assignment.

Remarks: This signature shall not participate in overload resolution unless `is_copy_constructible<E>::value and is_assignable<E &, E>::value`.

```cpp
template <class... Args>
void emplace(Args&&... args);
```

TODO: reword this operation

Effects: If `bool(*this)`, assigns the contained value `val` as if constructing an object of type `T` with the arguments `std::forward<Args>(args)`, otherwise destroys the contained value by calling `err->E::~E()` and initializes the contained value as if constructing an object of type `T` with the arguments `std::forward<Args>(args)`.

Postconditions: `bool(*this)`.
Exception Safety: If an exception is thrown during the call to \( T \)'s constructor or assignment, \(*\texttt{this}\) is disengaged, and the previous \( \text{val} \) (if any) has been destroyed.

Throws: Any exception thrown by the selected constructor or assignment of \( T \).

Remarks: This signature shall not participate in overload resolution unless \texttt{is\_constructible<T, Args&&...>::value}.

```cpp
template <class U, class... Args>
voidemplace(initializer_list<U> il, Args&&... args);
```

Effects: If \texttt{bool(*\texttt{this})}, assigns the contained value \( \text{val} \) as if constructing an object of type \( T \) with the arguments \( \text{il} \), \texttt{std::forward<Args>(args)}\..., otherwise destroys the contained value by calling \( \texttt{err->E::E()} \) and initializes the contained value as if constructing an object of type \( T \) with the arguments \( \text{il} \), \texttt{std::forward<Args>(args)}\...

Postconditions: \texttt{bool(*\texttt{this})}.

Exception Safety: If an exception is thrown during the call to \( T \)'s constructor, \! \texttt{bool(*\texttt{this})}, and the previous \( \text{val} \) (if any) has been destroyed.

Throws: Any exception thrown by the selected constructor of \( T \).

Remarks: The function shall not participate in overload resolution unless: \texttt{is\_constructible<T, initializer\_list<U>&, Args&&...>::value}.

X.Z.4.4 Swap [expected.object.swap]

```cpp
void swap(expected<T, E&, \texttt{rhs}) noexcept(/*see below*/);
```

Effects: If \texttt{bool(*\texttt{this}) and bool(\texttt{rhs})}, calls \texttt{swap(val, rhs.val)}\..., otherwise if \! \texttt{bool(*\texttt{this}) and bool(\texttt{rhs})}, calls \texttt{swap(err, rhs.err)}\..., otherwise if \texttt{bool(*\texttt{this}) and ! bool(\texttt{rhs})}, initializes a temporary variable \( e \) by direct-initialization with \texttt{std::move(\texttt{rhs.err})}\..., initializes the expected value of \*\texttt{this} by direct-initialization with \texttt{std::move(\texttt{rhs.err})} and swaps \texttt{has\_value} and \texttt{\texttt{has\_value}}\..., otherwise calls to \texttt{\texttt{rhs}swap(*\texttt{this})}.

Exception Safety: TODO: This must be worded.

Throws: Any exceptions that the expressions in the Effects clause throw.

Remarks: The expression inside noexcept is equivalent to:
\texttt{is\_nothrow\_move\_constructible<T>::value and noexcept(swap(declval<T&>(), declval<T&>()))} and \texttt{is\_nothrow\_move\_constructible<E>::value and noexcept(swap(declval<E&>(), declval<E&>()))}.

Remarks: The function shall not participate in overload resolution unless: \texttt{LValues of type T shall be Swappable, is\_move\_constructible<T>::value}, \texttt{LValues of type E shall be Swappable and is\_move\_constructible<T>::value}.

X.2.4.5 Observers [expected.object.observe]

```cpp
constexpr T\ const& operator->() const;
\end{verbatim}
\texttt{\texttt{T\& operator->()};}
```

Requires: \texttt{\texttt{bool(*\texttt{this})}}\).

Returns: \texttt{&val}.

Remarks: Unless \( T \) is a user-defined type with overloaded unary \texttt{operator\&}, the first function shall be a constexpr function.

```cpp
constexpr T\ const\& operator\*() const\&;
\end{verbatim}
\texttt{\texttt{T\& operator\*() \&};}
```

Requires: \texttt{\texttt{bool(*\texttt{this})}}\).

Returns: \texttt{val}.

Remarks: The first function shall be a constexpr function.

```cpp
constexpr T\&& operator\*() \&&;
\end{verbatim}
```

Requires: \texttt{\texttt{bool(*\texttt{this})}}\).

Returns: \texttt{move(val)}.

Remarks: This function shall be a constexpr function.
constexpr explicit operator bool() noexcept;

Returns: has_value.

Remarks: This function shall be a constexpr function.

constexpr T const& value() const;&
T& value() &;
constexpr T&& value() &&;

Returns: *val, if bool(*this).

Remarks: The first and third functions shall be constexpr functions.

constexpr E const& error() const;&
constexpr E& error() &;

Requires: ! bool(*this).

Returns: err.

Remarks: The first function shall be a constexpr function.

constexpr E&& error() &&;

Requires: ! bool(*this).

Returns: move(err).

Remarks: The first function shall be a constexpr function.

template<class Ex>
bool expected<T>::has_exception() const;

Returns: true if and only if ! bool(*this) and the stored exception is a base type of Ex.

constexpr unexpected_type<E> get_unexpected() const;

Requires: ! bool(*this).

Returns: make_unexpected(err).

template<class U>
constexpr T value_or(U&& v) const&;

Returns: bool(*this) ? **this : static_cast<T>(std::forward<U>(v)).

Exception Safety: If has_value and exception is thrown during the call to T’s constructor, the value of has_value and v remains unchanged and the state of val is determined by the exception safety guarantee of the selected constructor of T. Otherwise, when exception is thrown during the call to T’s constructor, the value of *this remains unchanged and the state of v is determined by the exception safety guarantee of the selected constructor of T.

Throws: Any exception thrown by the selected constructor of T.

Remarks: If both constructors of T which could be selected are constexpr constructors, this function shall be a constexpr function.

Remarks: The function shall not participate in overload resolution unless: is_copy_constructible<T>::value and is_convertible<U&&, T>::value.

template<class U>
T value_or(U& v) &&;

Returns: bool(*this) ? std::move(**this) : static_cast<T>(std::forward<U>(v)).

Exception Safety: If has_value and exception is thrown during the call to T’s constructor, the value of has_value and v remains unchanged and the state of val is determined by the exception safety guarantee of the T’s constructor.
Otherwise, when exception is thrown during the call to `T`'s constructor, the value of `*this` remains unchanged and the state of `v` is determined by the exception safety guarantee of the selected constructor of `T`.

**Throws:** Any exception thrown by the selected constructor of `T`.

**Remarks:** The function shall not participate in overload resolution unless: 

\[
\text{is_move_constructible<}T\text{>::value and is_convertible<U&&, T>::value.}
\]

```cpp
template<
class U, class E>
constexpr expected<U, E> expected<T, E>::unwrap() const;
```

**Returns:** If \[ bool(*this) \] then `**this` else `get_unexpected()`.

**Throws:** Any exception thrown by the selected constructor of `expected<U, E>`.

**Remarks:** The function shall not participate in overload resolution unless: 

\[
\text{is_move_constructible<}U\text{, }E\text{>::value.}
\]

```cpp
template<
class T, class E>
constexpr expected<T, E> expected<T, E>::unwrap() const;
```

**Returns:** `*this`.

**Throws:** Any exception thrown by the selected constructor of `expected<T, E>`.

**Remarks:** The function shall not participate in overload resolution unless: 

\[
T\text{ is not expected<}U\text{, }E\text{> and is_copy_constructible<}expected<T, E>\text{>::value.}
\]

```cpp
template<
class U, class E>
expected<T, E> expected<T, E>::catch_exception(F&& func);
```

**Effects:** if \[ has_exception<Ex>() \] call the continuation function `fuct` with the stored exception as parameter.

**Returns:** if \[ has_exception<Ex>() \] returns the result of the call continuation function `fuct` possibly wrapped on a `expected<T>`, otherwise, returns `*this`.

```cpp
template<
class Ex, class F>
'\text{see below}' map(F&& func)
```

**Returns:** if \[ bool(*this) \] returns `expected<\text{decltype(func(move(val)))}, E>{func(move(val))}`. otherwise, returns `get_unexpected()`.

```cpp
template<
class Ex, class F>
'\text{see below}' bind(F&& func)
```

**Returns:** if \[ bool(*this) \] returns `unwrap(expected<\text{decltype(func(move(val)))}, E>{func(move(val))})`. otherwise, returns `get_unexpected()`.

```cpp
template<
class Ex, class F>
'\text{see below}' then(F&& func);
```

**Returns:** `unwrap(expected<\text{decltype(func(move(*this)))}, E>{func(move(*this))})`.
template <class Ex, class F>
expected<T, E> catch_error(F && func);

Returns: if bool(*this) returns unwrap(expected<decaytype(func(val))), E(func(*this)), if ! bool(*this) returns the result of the call continuation function funct possibly wrapped on a expected<T, E>, otherwise, returns *this.

X.Z.6  expected for void [expected.object.void]

template <class E>
class expected<void, E> {
  public:
    typedef void value_type;
    typedef E error_type;
    template <class U>
      struct rebind {
        typedef expected<U, error_type> type;
      };
    // ??, constructors
    constexpr expected() noexcept;
    expected(const expected&);
    expected(expected& expected); noexcept (see below);
    expected(expected&& expected); noexcept (see below);
    constexpr expected(in_place_t);
    template <class Err>
      constexpr expected(unexpected_type<Err> const&);
    // ??, destructor
    -expected();
    // ??, assignment
    expected& operator=(const expected&);
    expected& operator=(expected& expected); noexcept (see below);
    void emplace();
    // ??, swap
    void swap(expected& expected); noexcept (see below);
    // ??, observers
    constexpr explicit operator bool() const noexcept;
    value() const;
    constexpr E const& error() const;
    constexpr E& error() &;
    constexpr unexpected_type<E> get_unexpected() const;
    template <typename Ex>
      bool has_exception() const;
    template 'see below' unwrap() const&;
    template 'see below' unwrap() &&;
    // ??, factories
    template <typename Ex, typename F>
      expected<void, E> catch_exception(F&& f);
    template <typename F>
      expected<decaytype(func()), E> map(F&& func);
    template <typename F>
      expected<void, E> catch_error(F&& f);
    template <typename F>
      'see below' then(F&& func);
  private:
    bool has_value; // exposition only
    union {
      unsigned char dummy; // exposition only
      error_type err; // exposition only
    };
};

TODO: Describe the functions

X.Z.7  unexpect  tag [expected.unexpect]

struct unexpect_t;
constexpr unexpect_t unexpect;

X.Z.8 Template Class  bad_expected_access  [expected.badexpectedaccess]
template <class E>
class bad_expected_access : public logic_error {

public:
    explicit bad_expected_access(E);
    constexpr error_type const& error() const;
    error_type& error();

};

The template class \texttt{bad\_expected\_access} defines the type of objects thrown as exceptions to report the situation where an attempt is made to access the value of a unexpected expected object.

\texttt{bad\_expected\_access::bad\_expected\_access(E e);}

\textbf{Effects:} Constructs an object of class \texttt{bad\_expected\_access} storing the parameter.

\texttt{constexpr E const\& \texttt{bad\_expected\_access::error()} const; \nE& \texttt{bad\_expected\_access::error();}}

\textbf{Returns:} The stored error.

\textbf{Remarks:} The first function shall be a constexpr function.

X.Z.8 Expected Relational operators [expected.relational.op]

\textbf{template <class T, class E>}
\texttt{constexpr bool operator==(const expected\_T, E\& x, const expected\_T, E\& y);} 

\textbf{Requires:} T and unexpected\_type\<E\> shall meet the requirements of EqualityComparable.

\textbf{Returns:} If \texttt{bool(x) != bool(y)}, \texttt{false} ; otherwise if \texttt{bool(x) == false} , \texttt{x.get\_unexpected()} == \texttt{y.get\_unexpected()} ; otherwise \texttt{*x == *y}.

\textbf{Remarks:} Specializations of this function template, for which \texttt{*x == *y} and \texttt{x.get\_unexpected()} == \texttt{y.get\_unexpected()} are core constant expression, shall be constexpr functions.

\textbf{template <class T, class E>}
\texttt{constexpr bool operator!=(const expected\_T, E\& x, const expected\_T, E\& y);} 

\textbf{Returns:} \texttt{!(x == y)}. 

\textbf{template <class T, class E>}
\texttt{constexpr bool operator<(const expected\_T, E\& x, const expected\_T, E\& y);} 

\textbf{Requires:} \texttt{*x < *y} shall be well-formed and its result shall be convertible to \texttt{bool}.

\textbf{Returns:} \texttt{!(y < x)}. If \texttt{y < x} , \texttt{x.get\_unexpected()} < \texttt{y.get\_unexpected()} ; otherwise \texttt{*x < *y}.

\textbf{Remarks:} Specializations of this function template, for which \texttt{*x < *y} and \texttt{x.get\_unexpected()} < \texttt{y.get\_unexpected()} are core constant expression, shall be constexpr functions.

X.Z.9 Comparison with \texttt{T} [expected.comparison_T]
template <class T, class E> constexpr bool operator==(const T& x, const T& v);
template <class T, class E> constexpr bool operator==(const E& x, const T& v);


template <class T, class E> constexpr bool operator!=(const T& x, const T& v);
template <class T, class E> constexpr bool operator!=(const E& x, const T& v);

Returns: !(x == v).

template <class T, class E> constexpr bool operator<(const T& x, const T& v);


template <class T, class E> constexpr bool operator<=(const T& x, const T& v);

Returns: boolxpr ? v < *x : false.

X.Z.10 Comparison with unexpected_type<E> [expected.comparisonunexpectedE]

template <class T, class E> constexpr bool operator==(const T& x, const unexpected_type<E>& v);
template <class T, class E> constexpr bool operator==(const unexpected_type<E>& x, const T& v);


template <class T, class E> constexpr bool operator!=(const T& x, const unexpected_type<E>& v);
template <class T, class E> constexpr bool operator!=(const unexpected_type<E>& x, const T& v);

Returns: !(x == v).

template <class T, class E> constexpr bool operator<(const T& x, const unexpected_type<E>& y);

Returns: y < x.

template <class T, class E> constexpr bool operator<=(const T& x, const unexpected_type<E>& y);

Returns: y < x.
template <class T, class E> constexpr bool operator >= (const expected<T, E>& x, constunexpected_type<E>& y);

Returns: !(x < y).

X.Z.11 Specialized algorithms [expected.specalg]

template <class T, class E>
void swap (expected<T, E>& x, expected<T, E>& y) noexcept (noexcept (x.swap(y)));

Effects: calls x.swap(y).

X.Z.12 Expected Factories [expected.factories]

template <class T, class E>
constexpr expected<typename decay<T>::type> make_expected (T&& v);

Returns: expected<typename decay<T>::type>(std::forward<T>(v)).

template <class T, class E>
expected<T, exception_ptr> make_expected_from_exception (std::exception_ptr v);

Returns: expected<T, exception_ptr>(unexpected_type<E>(std::forward<E>(v))).

template <class T, class E>
expected<T, decay_t<E>> make_expected_from_error (E&& e);

Returns: expected<T, decay_t<E>>(make_unexpected(e)).

template <class T, class E, class U>
expected<T, E> make_expected_from_error (U&& u);

Returns: expected<T, E>(make_unexpected(E{forward<U>(u)})).

template <class T, class E>
expected<T, exception_ptr> make_expected_from_current_exception ();

Returns: expected<T, exception_ptr>(make_unexpected_from_current_exception()).

template <class F>
constexpr typename expected_result_of<F>::type make_expected_from_call (F funct);

Equivalent to:

try {
  return make_expected(funct());
} catch (...) {
  return make_unexpected_from_current_exception();
}

X.Z.13 Hash support [expected.hash]

template <class T, class E>
struct hash<expected<T, E>>;

Requires: The template specializations hash<T> and hash<E> (if E is not exception_ptr) shall meet the requirements of class template hash (Z.X.Y). The template specialization hash<expected<T, E>> shall meet the requirements of class template hash. For an object e of type expected<T, E>, if bool(e).hash<expected<T, E>>()(e) shall evaluate to a combination of the hashing true and hash<T>()(e.error()); otherwise it evaluates to an unspecified value if E is exception_ptr or a combination of hashing false and hash<E>()(e.error()).

template <class E>
struct hash<expected<void, E>>;

Requires: The template specialization hash<E> (if E is not exception_ptr) shall meet the requirements of class template hash (Z.X.Y). The template specialization
hash<expected<void,E>> shall meet the requirements of class template hash. For an object e of type expected<void,E>, if bool(e), hash<expected<void,E>>()(e) shall evaluate to the hashing true; otherwise it evaluates to an unspecified value if E is exception_ptr, or a combination of hashing false and hash<E>()(e.error()).

X.Z.14 expected as a meta-fuction [expected.object.meta]

```
template <class E>
class expected<holder, E> {
public:
    template <class T>
    using apply = expected<T, E>;
};
```

### Implementability

This proposal can be implemented as pure library extension, without any compiler magic support, in C++14. An almost full reference implementation of this proposal can be found at TBoost.Expected.

### Future Work

### Allocator support

As optional<T>, expected<T,E> does not allocate memory. So it can do without allocators. However, it can be useful in compound types like:

```
typedef vector<expected<vector<int, MyAlloc>, error>, MyAlloc> MyVec;
MyVec v{ v2, MyAlloc() };
```

One could expect that the allocator argument is forwarded in this constructor call to the nested vectors that use the same allocator. Allocator support would enable this. `std::tuple` offers this functionality.

### Variadic expected

A typical case could combine expected and variant expected<T,variant<E1, ..., En>>. We could extend expected to take a variadic number of errors expected<T, E1, ..., En> in order to provide an adapted interface.

### Acknowledgements

We are very grateful to Andrei Alexandrescu for his talk, which was the origin of this work. We thanks also to every one that has contributed to the Haskell either monad, as either's interface was a source of inspiration.

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Vicente thanks personally Evgeny Panasyuk and Johannes Kaphammer for their remarks on the DO-expression. Pierre thanks the IRCAM and Carlos Agon who allowed him to work on this proposal during its internship.

### References


- [ERR](https://github.com/psiha/err)

  err - yet another take on C++ error handling, 2015.

- [N3527](http://www.open-std.org/jtc1/sc22/wg21/docs/papers/2013/n3527.html)

  Fernando Cacciola and Andrzej Krzemieński. N3527 - A proposal to add a utility class to represent optional objects (Revision 3), March 2013.

- [N3672](http://www.open-std.org/jtc1/sc22/wg21/docs/papers/2013/n3672.html)

  Fernando Cacciola and Andrzej Krzemieński. N3672 - A proposal to add a utility class to represent optional objects (Revision 4), June 2013.

- [N3793](http://www.open-std.org/jtc1/sc22/wg21/docs/papers/2013/n3793.html)

  Fernando Cacciola and Andrzej Krzemieński. N3793 - A proposal to add a utility class to represent optional objects (Revision 5), October 2013.
In the use-cases section, we present expected as a better way to handle errors than error codes. However the current syntax using lambda and chaining monadic operations (such as `map`) can be tedious to use. We propose different solutions to overcome this problem, since the solutions are more general than the scope of this proposal we discuss them in this appendix.

A first solution that do not require change in the language is the use of variadic monadic operation. For example using a variadic free function `map`, we can write the function as following:

```cpp
expected<int> f(int i, int j, int k)
{
    return map(plus,
        safe_divide(i, k),
        safe_divide(j, k));
}
```

This is most readable than the member map function and the use of lambda. However it suffers from two major deficiencies:

- **Eager evaluation**: All arguments are evaluated even if the first fails.
- **Unordered evaluation**: We cannot control the order of evaluation thus it presupposed the function to have no side effects.

Considering these two problems we consider a possible C++ language extension: a `do-notation` similar to the one provided by Haskell [Do-Notation]. As with the variadic `map` function, it is not limited to expected but could work with any kind of monad [Monad]. Next follows the grammar:

**Appendix I - Language support**

**Better support for monad**

In the use-cases section, we present expected as a better way to handle errors than error codes. However the current syntax using lambda and chaining monadic operations (such as `map`) can be tedious to use. We propose different solutions to overcome this problem, since the solutions are more general than the scope of this proposal we discuss them in this appendix.

A first solution that do not require change in the language is the use of variadic monadic operation. For example using a variadic free function `map`, we can write the function as following:

```cpp
expected<int> f(int i, int j, int k)
{
    return map(plus,
        safe_divide(i, k),
        safe_divide(j, k));
}
```
expression ::= ...  
  | do-expression

do-expression ::= do-initialization ':=' expression

do-initialization ::= type var '<-' expression

The previous function could be rewritten as:
```c
expected<int, error_condition> f2(int i, int j, int k) {
  return ( auto s1 <- safe_divide(i, k) :
          auto s2 <- safe_divide(j, k) : s1 + s2
```

This syntax is far easier to read and to understand. A lazy evaluation of statement is possible and the order is well-defined. Nevertheless, it can be considered as a syntactic sugar for bind. We give a syntactic transformation following:

```c
[[do-expression]] =
  bind(expression, [=](type var) {
    return [[do-expression-or-expression]];  
  });
```

The transformed code of the previous function is:
```c
expected<int, error_condition> f2(int i, int j, int k)
{
  return bind(safe_divide(i, k), [=](auto s1) {
    return bind(safe_divide(j, k), [=](auto s2) {
      return s1 + s2;
    });
  });
}
```

This would give the exact same results as the previous version. However, the function f2 is much simpler and clearer than f because it doesn’t have to explicitly handle any of the error cases. When an error case occurs, it is returned as the result of the function, but if not, the correct result of a subexpression is bound to a name (s1 or s2), and that result can be used in later parts of the computation. The code is a lot simpler to write.

The more complicated the error-handling function, the more important this will be. But, the standard doesn’t have this DO expression yet. Waiting for a do-statement the user could define some macros [DO-Macro] and define f2 as
```c
expected<int> f2(int i, int j, int k)
{
  return DO(
    ( s1, safe_divide(i, k) )
    ( s2, safe_divide(j, k) )
  s1 + s2
  );
}
```

In the case of expected and optional, and similarly to the proposed await keyword we could use an expect keyword (it returns the unexpected value if the expected is not valued):
```c
expected<int> f2(int i, int j, int k)
{
  auto s1 = expect safe_divide(i, k);
  auto s2 = expect safe_divide(j, k);
  return s1 + s2;
}
```

We could even nest the expect expression as in
```c
expected<int> f2(int i, int j, int k)
{
  return expect safe_divide(i, k) + expect safe_divide(j, k);
}
```

but these expressions cannot be computed concurrently, we need a specific order of evaluation.

Waiting for this, the user would define a macro [EXPECT-Macro]
Expected might be configurable through a trait expected_traits.

The first variation point is the behavior of value() when expected<T,E> contains an error. The current strategy throw a bad_expected_access exception (or the contained exception if the type is expected<T,exception_ptr>) but it might not be satisfactory for every error type. For example, some might want to encapsulate an error_condition into a specific exception. Or in debug mode, they might want to use an assert call.

The other variation point is the behavior triggered when the continuation argument of bind or map throws an exception. If the exception thrown is system_error and the error type is error_code, we might want to store the error carried by the exception. Without more discussion, let’s show how we could customize expected<T, E>, consider the following exception-oriented function:

Appendix II - Alternative designs

A Configurable Expected

Expected might be configurable through a trait expected_traits.

The first variation point is the behavior of value() when expected<T,E> contains an error. The current strategy throw a bad_expected_access exception (or the contained exception if the type is expected<T,exception_ptr>) but it might not be satisfactory for every error type. For example, some might want to encapsulate an error_condition into a specific exception. Or in debug mode, they might want to use an assert call.

The other variation point is the behavior triggered when the continuation argument of bind or map throws an exception. If the exception thrown is system_error and the error type is error_code, we might want to store the error carried by the exception. Without more discussion, let’s show how we could customize expected<T, E>, consider the following exception-oriented function:
class error_cond : public std::exception {
  // Implementation similar to system_error but for error_condition here.
};
int safe_divide(int i, int j) {
  if (j == 0)
    throw error_cond(error_condition(arithmetic_errc::divide_by_zero));
  return i/j;
}

Imagine \( j \) encapsulated into an \texttt{expected} , you will call map with \texttt{safe_divide} as the continuation. Let’s see what it looks like:

\[
\text{expected-int, error_condition} \ f(int \ i, \ \text{const expected-int, error_condition} & j) \\
\{
  return j.map(bind(safe_divide, i, _1));
\}
\]

If we specialize \texttt{expected_traits} for \texttt{error_condition} , we can achieve the expected behavior:

\[
\text{template <class T> struct expected_traits<expected<T, error_condition>>} \\
\{ \\
  static expected<T, error_condition> catch_exception(exception_ptr e) \\
  { \\
    try{ \\
      rethrow_exception(e); \\
    } catch(const error_cond & e) { \\
      return make_unexpected(e.code());
    }
  }
  static void bad_access(const error_type &e) \\
  { \\
    throw error_cond(e);
  }
\};
\]

The semantics of \texttt{catch_exception} is to re-throw the current exception and catch only the exceptions we are interested in. The default behavior lets flight the exception thrown by the continuation. We created a bridge between an \texttt{error_condition} and the \texttt{error_cond} exception.

**Which exception throw when the user try to get the expected value but there is none?**

It has been suggested to let the user decide the exception that would be throw when the user try to get the expected value but there is none, as third parameter.

While there is no major complexity doing it, as it just needs a third parameter that could default to the appropriated class,

\[
\text{template <class T, class Error, class Exception = bad_expected_access> struct expected;} \\
\]

The authors consider that this is not really needed and that this parameter should not really be part of the type.

The user could use \texttt{value_or_throw()}

\[
\text{expected-int, std::error_code} \ f(); \\
\text{expected-int, std::error_code} \ e = f(); \\
\text{auto i} = e.value_or_throw<std::system_error>();
\]

where

\[
\text{template <Exception, class T, class E> constexpr value_type value_or_throw(expected-T,E&, e) const&} \\
\{ \\
  return *this \\
  ? move(**this) \\
  : throw Exception(e.error());
\}
\]

A function like this one could be added to the standard, but this proposal doesn’t request it. The user can also wrap the proposed class in its own expected class.
template <class T, class Error=std::error_code, class Exception=std::system_error>
struct MyExpected {
    expected<T, E> v;
    MyExpected(expected<T, E> v) : v(v) {
        T value() {
            if (e) return v.value();
            else throw Exception(v.error());
        }
    }
    ...;
};

and use it as

    expected<int, std::error_code> f();
    MyExpected<int> e = f();
    auto i = e.value(); // std::system_error throw if not valid

A class like this one could be added to the standard, but this proposal doesn't request it.

An alternative could be to add a specialization on an error class that gives the storage and the exception to thrown.

template <class Error, class Exception>
struct error_exception {
    typedef Error error_type;
    typedef Exception exception_type;
};

that could be used as follows

    expected<T, std::error_exception::error_code, std::system_error> e = make_unexpected(err);
    e.value(); // will throw std::system_error(err);

A class like this one could be added to the standard, but this proposal doesn't request it.

About expected

It has been suggested also to extend the design into something that contains

- a ?
- an ErrorCode throw using Exception, or
- a exception_ptr

Again there is no major difficulty to implement it, but instead of having one variation point we have two, that is, is there a value, and if not, if is there an exception_ptr. While this would need only an extra test on the exceptional case, the authors think that it is not worth doing it as all the copy/move/swap operations would be less efficient.

Appendix III - Related types

Variant

expected<T, E> can be seen as a specialization of boost::variant<unexpected_type<E>, T> which gives a specific intent to its first parameter, that is, it represents the type of the expected contained value. This specificity allows to provide a pointer like interface, as it is the case for std::experimental::optional<T>. Even if the standard included a class variant<T, E>, the interface provided by expected<T, E> is more specific and closer to what the user could expect as the result type of a function. In addition, expected<T, E> doesn't intend to be used to define recursive data as boost::variant<> does.

The following table presents a brief comparison between boost::variant<unexpected_type<E>, T> and expected<T, E>. 

We can see `expected<T,E>` as an `std::experimental::optional<T>` that collapse all the values of `E` to `nullopt`.

We can convert an `expected<T,E>` to an `optional<T>` with the possible loss of information.

```cpp
template <class T> optional<T> make_optional(expected<T,E> v) {
    if (v) return make_optional(*v);
    else nullopt;
}
```

We can convert an `optional<T>` to an `expected<T,E>` without knowledge of the root cause. We consider that `E()` is equal to `nullopt` since it shouldn’t bring more informations (however it depends on the underlying error — we considered `exception_ptr` and `error_condition`).

```cpp
template <class T, class E> expected<T,E> make_expected(optional<T> v) {
    if (v) return make_expected(*v);
    else make_unexpected(E());
}
```

## Promise and Future

We can see `expected<T>` as an always ready `future<T>`. While `promise<> / future<>` focuses on inter-thread asynchronous communication, `expected<E,T>` focus on eager and synchronous computations. We can move a ready `future<T>` to an `expected<T>` with no loss of information.

```cpp
template <class T> expected<T> make_expected(future<T>&& f) {
    assert(f.ready() && "future not ready");
    try {
        return f.get();
    } catch (...) {
        return make_unexpected_from_exception();
    }
}
```

We could also create a `future<T>` from an `expected<T>`

```cpp
template <class T> future<T> make_future(expected<T> e) {
    if (e)
        return make_ready_future(*e);
    else
        return make_exceptional_future<T>(e.error());
}
```

## Comparison between optional, expected and future
The following table presents a brief comparison between `optional<T>`, `expected<T,E>` and `promise<T> / future<T>`.

<table>
<thead>
<tr>
<th></th>
<th>optional</th>
<th>expected</th>
<th>promise/future</th>
</tr>
</thead>
<tbody>
<tr>
<td>specific null value</td>
<td>yes</td>
<td>no</td>
<td>non</td>
</tr>
<tr>
<td>relational operators</td>
<td>yes</td>
<td>yes</td>
<td>no</td>
</tr>
<tr>
<td>swap</td>
<td>yes</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>factories</td>
<td>make_option / nullopt</td>
<td>make_expected / make_unexpected</td>
<td>make_ready_future / (make_exceptional, see [P0159R0])</td>
</tr>
<tr>
<td>hash</td>
<td>yes</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>value_type</td>
<td>yes</td>
<td>yes</td>
<td>no / (yes, see [N3865]).</td>
</tr>
<tr>
<td>default constructor</td>
<td>yes</td>
<td>yes (if T is default constructible)</td>
<td>yes</td>
</tr>
<tr>
<td>allocators</td>
<td>no</td>
<td>no</td>
<td>yes</td>
</tr>
<tr>
<td>emplace</td>
<td>yes</td>
<td>yes</td>
<td>no</td>
</tr>
<tr>
<td>bool conversion</td>
<td>yes</td>
<td>yes</td>
<td>no</td>
</tr>
<tr>
<td>state</td>
<td>bool()</td>
<td>bool() / valid</td>
<td>valid / ready / (has_value, see [N3865])</td>
</tr>
<tr>
<td>observers</td>
<td>pointer-like / value / value_or</td>
<td>pointer-like / value / error / value_or</td>
<td>get / (get_exception_ptr, see [N3865])</td>
</tr>
<tr>
<td>visitation</td>
<td>no</td>
<td>map / bind / then / catch_error</td>
<td>then / (next/recover see [N3865])</td>
</tr>
<tr>
<td>grouping</td>
<td>n/a</td>
<td>n/a</td>
<td>when_all / when_any</td>
</tr>
</tbody>
</table>