

Document number: P0075R0  
 Date: 2015-9-25  
 Project: Programming Language C++, Library Working Group  
 Reply to: Arch D. Robison <arch.robison@intel.com>  
 Pablo Halpern <pablo.g.halpern@intel.com >  
 Robert Geva <robert.geva@intel.com>  
 Clark Nelson <clark.nelson@intel.com>

## Template Library for Index-Based Loops

### 1 Introduction

Indexed-based loops are well established in programming languages. Though C++ has language-level support for sequential forms of these loops, it has none for parallel or parallel-vector forms. This paper proposes support for all aforementioned forms as a pure-library extension. Our proposal is pure-library extension of the Parallelism TS, and adds support for indexed-based loops with reduction and induction variables.

The proposal adds the following new function templates to the Parallelism TS:

- **for\_loop** and **for\_loop\_strided** implement loop functionality over a range specified by integral or iterator bounds. For the iterator case, it resembles `for_each` from the Parallelism TS, but leaves to the programmer when and if to dereference the iterator.
- **reduction** provides a flexible way to specify reductions in conjunction with `for_loop`.
- **reduction\_plus**, **reduction\_multiplies**, ... etc. creating reduction descriptors for common cases such as addition, multiplication, etc.
- **induction** provides a flexible way to specify indices or iterators that vary linearly with the primary index of the loop.

Here is a short example:

```
void saxpy_ref(int n, float a, float x[], float y[]) {
    for_loop( seq, 0, n, [&](int i) {
        y[i] += a*x[i];
    });
}
```

The call to `for_loop` is equivalent to:

```
void saxpy_ref(int n, float a, float x[], float y[]) {
    for( int i=0; i<n; ++i )
        y[i] += a*x[i];
}
```

The loop can be parallelized by replacing `seq` with `par`. Our library interface permits the “loop index” to have integral type or be a random-access iterator. As with the current Parallelism TS, the iterator case

does not require a random-access iterator. For example, `for_loop` enables the following general implementation of `for_each` from the Parallelism TS.

```
template<class ExecutionPolicy,
class InputIterator, class Function>
void for_each(ExecutionPolicy&& exec, InputIterator first, InputIterator last,
             Function f) {
    for_loop( exec, first, last, [&](InputIterator i){f(*i);} );
}
```

When `exec` is not `sequential_execution_policy`, random-access iterators may yield better performance, because unaggressive implementations are likely to fall back to using a serial loop for other kinds of iterators.

## 1.1 Strided Loops

Our proposal also adds a function template for strided loops. Though these can be expressed from unit-stride loops and mathematical machinations, we think code is clearer when loops can be expressed in natural strided form. To alleviate template overload trickiness and potential hazards, we the function template for strided loops has a different name. The situation is somewhat akin to the motivations for giving `for_each` and `for_each_n` different names.

The stride parameter follows the second bound on the index space. The `stride=3` example below sets `c[10]`, `c[13]`, `c[16]`, `c[19]` to true.

```
for_loop_strided( par, 10, 20, 3, [&](int k) {
    c[k] = true;
});
```

Negative strides are allowed. The following sets the same elements of `c` to true as the previous example.

```
for_loop_strided( par, 19, 9, -3, [&](int k) {
    c[k] = true;
});
```

## 1.2 Reductions

The `for_loop` template also allows specification of one more reduction variables, with a syntax inspired by OpenMP, but done with a pure library approach. Here is an example:

```
float dot_saxpy(int n, float a, float x[], float y[]) {
    float s = 0;
    for_loop( par, 0, n,
             reduction(s,0.0f,std::plus<float>()),
             [&](int i, float& s_) {
                 y[i] += a*x[i];
                 s_ += y[i]*y[i];
             });
    return s;
}
```

Here, `reduction` is a function that returns an implementation-specified *reduction object* that specifies three things:

- a reduction lvalue  $s$
- the identity value for the reduction operation
- the reduction operation

In the lambda expression,  $i$  is a value of the loop index, and  $s_*$  is a reference to a private partial sum. There is one such reference for each reduction argument to `for_loop`, and association is positional. (We suspect that in practice, most programmers will name the local reference just  $s$ .) The example is equivalent, except with more relaxed sequencing and reduction order, to the following serial code:

```
float serial_dot_saxpy (int n, float a, float x[], float y[]) {
    float s = 0;
    for( int i=0; i<n; ++i ) {
        y[i] += a*x[i];
        s += y[i]*y[i];
    }
    return s;
}
```

For convenience, we supply shorthand functions for common reductions. For example:

```
reduction_plus(s)
```

is equivalent to:

```
reduction(s,0.0f,std::plus<float>())
```

### 1.3 Inductions (Linear Variables)

The for-loop template also allows specification of induction variables, using a scheme somewhat similar to that for reduction variables. Here is an example with three induction variables:

```
float* zipper(int n, float* x, float *y, float *z) {
    for_loop( par, 0, n,
             induction(x),
             induction(y),
             induction(z,2),
             [&](int i, float* x_, float* y_, float* z_) {
                 *z_++ = *x_++;
                 *z_++ = *y_++;
             });
    return z;
}
```

Here `induction` is a function that returns an implementation-specified type that specifies two things:

- a reference to an induction lvalue (e.g.  $x$ )
- a optional stride for that lvalue. Here the stride is implicitly 1 for  $x$  and  $y$ , and explicitly 2 for  $z$ .

In the lambda expression,  $i$  is a value of the loop index, and  $x_*$ ,  $y_*$ ,  $z_*$  are initialized with  $x+i$ ,  $y+i$ , and  $z+2*i$  respectively. As with reduction arguments, association is positional. A function can have both reduction and induction arguments. When the `for_loop` finishes,  $x$ ,  $y$ ,  $z$  are set to the same live-out

values as if the loop had been written sequentially. For example, the following serial code returns the same value as the previous example:

```
float* zipper(int n, float* x, float *y, float *z) {
    for( int i=0; i<n; ++i ) {
        *z++ = *x++;
        *z++ = *y++;
    }
    return z;
}
```

## 2 Alternative Designs

It is possible to leave induction out and rely on users to write the equivalent math. However, doing so complicates parallelizing codes. We note that OpenMP has linear clauses for similar reason.

The current Parallel STL has support for reductions. However, these are tightly tied to specific algorithms and require “tuple-fying” values (and defining reduction operations on the tuples) for code that needs to perform more than one reduction. Our approach brings the flexibility that OpenMP users have enjoyed from the start.

## 3 C++ Proposed Wording

The proposed edits are with respect to the current Parallelism TS.

### Reduction Support for `for_loop` [\[Addition to Non-Numeric Parallel Algorithms\]](#)

Reduction objects add a flexible reduction capability to `std::for_loop`. Reduction objects have implementation-specified types, and are created by the function template `reduction`.

```
namespace std {
namespace experimental {
namespace parallel {
inline namespace v2 {

// General form for reduction
template<typename T, typename Op>
implementation-specified reduction( T& var, T&& identity, Op&& op );

// Shorthand for plus reduction
template<typename T>
implementation-specified reduction_plus( T& var );

// Shorthand for multiplies reduction
template<typename T>
implementation-specified reduction_multiplies( T& var );

// Shorthand for bit_and reduction
template<typename J>
implementation-specified reduction_bit_and( J& var );

// Shorthand for bit_or reduction
template<typename J>
```

```

implementation-specified reduction_bit_or( J& var );

// Shorthand for bit_xor reduction
template<typename J>
implementation-specified reduction_bit_xor( J& var );

// Shorthand for min reduction
template<typename T>
implementation-specified reduction_min( T& var );

// Shorthand for max reduction
template<typename T>
implementation-specified reduction_max( T& var );

}}}}

```

Each function returns a “reduction object” that specifies a *var*, an *identity* value for the reduction, and a *reduction-op*. See description of `for_loop` for how these are used. The implicit *identity* and *reduction-op* are as follows:

<i>function</i>	<i>identity</i>	<i>reduction-op</i>
<code>reduction_plus</code>	<code>T()</code>	<code>x+y</code>
<code>reduction_multiplies</code>	<code>T(1)</code>	<code>x*y</code>
<code>reduction_bit_and</code>	<code>~(T())</code>	<code>x&amp;y</code>
<code>reduction_bit_or</code>	<code>T()</code>	<code>x y</code>
<code>reduction_bit_xor</code>	<code>T()</code>	<code>x^y</code>
<code>reduction_min</code>	<code>std::numeric_limits&lt;T&gt;::max()</code>	<code>std::min(x,y)</code>
<code>reduction_max</code>	<code>std::numeric_limits&lt;T&gt;::lowest()</code>	<code>std::max(x,y)</code>

[Example:

The following code updates each element of *y* and sets *s* to the sum of the squares.

```

float s = 0;
for_loop( vec, 0, n,
    reduction(s, std::plus<float>()),
    [&](int i, float& t) {
        y[i] += a*x[i];
        t += y[i]*y[i];
    }
});

```

--end example]

## Induction objects

Induction objects add a flexible capability to specify secondary index variables to `std::for_loop`.

```

namespace std {
namespace experimental {
namespace parallel {

```

```

inline namespace v2 {

template<typename T>
implementation-specified induction( T& var );

template<typename T, typename S>
implementation-specified induction( T& var, S stride );

}}}}

```

Each function returns an “induction object” that specifies a *var*, and optionally a *stride*. See description of `for_loop` for how these are used.

### For loop [\[Addition to Non-Numeric Parallel Algorithms\]](#)

```

namespace std {
namespace experimental {
namespace parallel {
inline namespace v2 {
template<typename Policy, typename I, typename... Rest>
void for_loop ( Policy&& policy, I first, I last, Rest&&... rest );

template<typename Policy, typename I, typename S, typename... Rest>
void for_loop_strided( Policy&& policy, I first, I last, S stride, Rest&&... rest );
}}}}

```

*Requires:* The parameter pack `rest` shall have at least one element; the last element shall be an invocable object, *f*, with an argument list composed as described below. Every other element of the parameter pack shall be the result of invoking an instance of one of the reduction or induction function templates.

*stride* shall be non-zero.

The type of *f* shall meet the requirements of `MoveConstructible` if `Policy` is `sequential_execution_policy`, otherwise it shall meet the requirements of `CopyConstructible`. For `for_loop`, *I* shall be an integral type or meet the requirements of an input iterator. For `for_loop_strided`, *S* shall be an integral type and:

- *I* shall be an integral type or
- *I* shall meet the requirements of an input iterator if *stride* is positive or
- *I* shall meet the requirements of a bidirectional iterator if *stride* is negative.

*Effects:* Applies *f* to a sequence of argument lists. For `for_loop`, the value of the first argument is successive values in the range `[first,last)`. For `for_loop_strided` with a positive *stride*, the value of the first argument is successive values in the range `[first,last)` starting with *first* and advancing by *stride*. For `for_loop_strided` with negative *stride*, the value of the first argument is successive values in the range `(last,first]`, starting with *first* and advancing (backwards) by *stride*.

Each argument list contains an additional iteration-local value corresponding, by positional order, to each parameter in `rest`, except the last.

For a reduction object, its iteration-local value is a reference to a temporary of its reduction type. Each temporary is copy-constructed from the reduction's *identity* value. The reduction object's *var* is updated with the result of applying *reduction-op* to *var* and any temporaries that were generated. When these updates occur is implementation specific.

For induction object with no stride, its iteration-local value for the *i*th iteration is *original + i* and its *final value* is *var + n*, where *n* is the number of times that *f* was applied. For an induction object with a *stride*, its iteration-local value for the *i*th iteration is *original + i\*stride* and its *final value* is *var + n\*stride*. The *var* is set to the *final value*. All applications of *f* are sequenced before *var* is set.

If *Policy* is `sequential_execution_policy`, the application shall start with the first argument list in the sequence and proceed to the last one.

*Complexity:* `for_loop` applies *f* exactly `last-first` times. `for_loop_strided` applies *f* exactly  $\lfloor (\text{last}-\text{first}-1)/\text{stride} \rfloor + 1$  if *stride* is positive, and  $\lfloor (\text{first}-\text{last}-1)/\text{stride} \rfloor + 1$  times if *stride* is negative.

*Remarks:* If *f* returns a result, the result is ignored.

[*Note:* When *i* and *j* are iterators, `for_loop(policy, i, j, f)` differs from `for_each(policy, i, j, f)` in that the latter dereferences elements in `[first,last)` --*note*]