

PRODUCTS GROUP

Java Study Group Oct 26-27

Roger Golliver Floating-point Center of Expertise roger.a.golliver@intel.com

Intel/FP-COE

Some of Java's Problematic Floating-point Requirements

- Java requires orthogonal IEEE default behavior
 - No unmasked exceptions
 - No IEEE flags, (but this should and could be easily corrected)
 - Bit for bit exact results, i.e. no double rounding errors from extended register's extra range and/or precision
- First, this ignores the installed base of PC using extended IEEE floating-point.
- Second, I know of no conforming multiply/divide implementations on extended architectures.

Sun is currently trying to address this Floating-point issue

- Sun is using its "open" PAS process
- Draft proposal published on the Internet
- Public comments were request by Sept 15, 1998
- One organized public response was from the JavaGrande group
- Sun's response to the public comments has not been made public yet.

The proposal will likely contain two floating-point modes

- The default mode will be tolerant of different IEEE/ANSI Std. 754-1985 conforming hardware implementations
- The second mode will a "strict" orthogonal mode
 - a new keyword will be introduced "strict"
 - it will conform to Java's current bit-for-bit exact FP results

Sun's draft proposal's algorithm to get strict multiply results on extended architectures

fld	qword ptr [dx]	/* dz = dx * dy */
fclex		/* clear flags */
fmul	qword ptr [dy]	/* 53-bits of sign., 15-bits of exp. */
fstsw	word ptr [sw]	/* rounded-up in Cl
		and sticky in Precision(Inexact) */
fst	qword ptr [dtmp]	/* 53-bits of significand, */
fstsw	ax	/* and 11-bits of exponent */
and	ax,0x30	/* Precision/Inexact AND Underflow */
xor	ax,0x30	/* set after fmul and store? */
jne	skip	/* if not then okay, continue */
jsr	fix_up	/* fix-up will use [sw] and top of x87 to
		round and clamp as required by STD Java */
in:		

skip:

fstp qword ptr [dz]

New algorithm to perform a strict multiply on an extended IEEE architectures

- precision control set to
 53-bits
 x_de = x_d
 x_de *= 2.0^(Emax_d-Emax_de)
 y_de = y_d
 x_de = x_de * y_de
 x_de *= 2.0^(Emax_de-Emax_d)
- Assume IA-32[™] style architecture exact, promotion exact will scale down exact, promotion will underflow correctly (denormalize) if tiny exact will scale up will overflow correctly
- will overflow correctly if huge
- Emax_de=0x7FFE-0x3FFF=0x3FFF
 Emax_d =0x7FE-0x3FF=0x3FF
- $Emax_de-Emax_d = 0x3C00$ $Emax_d-Emax_de = -0x3C00$

z d = x de

6

•Advantages of the new algorithm

- No expensive serializing operations on the control and status words.
- It allows for optimizations which hide the latency of floating-point operations.
- Its cost is only two more multiples from what current JVM's are probably doing.

7

Conclusion and ways to follow up

- Questions on the new algorithm can be directed to me at roger.a.golliver@intel.com.
- JavaGrande is group interested in other issues related to using Java for Scientific and Engineering applications.
 - you can visit them at: http://www.javagrande.org
 - JavaGrande will have a "Birds-of-a-Feather" session at the SC98 (Supercomputing98) in November.